# Relation Between Randić Index and Average Distance of Trees ${ }^{1}$ 

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#### Abstract

The Randić index $R(G)$ of a graph $G$ is the sum of weights $(\operatorname{deg}(u) \operatorname{deg}(v))^{-0.5}$ over all edges $u v$ of $G$, where $\operatorname{deg}(v)$ denotes the degree of a vertex $v$. We prove that for any tree $T$ with $n_{1}$ leaves $R(T) \geq \operatorname{ad}(T)+\max \left(0, \sqrt{n_{1}}-2\right)$, where $\operatorname{ad}(T)$ is the average distance between vertices of $T$. As a consequence we resolve the conjecture $R(G) \geq \operatorname{ad}(G)$ given by Fajtlowicz in 1988 for the case when $G$ is a tree.


## 1 Introduction

In chemical graph theory topological indices belong to the set of molecular descriptors that are calculated based on the molecular graph of a chemical compound. In 1975 Milan Randić [11] introduced a topological connectivity index $R(G)$ of a graph $G$ defined as the sum of weights $(\operatorname{deg}(u) \operatorname{deg}(v))^{-0.5}$ over all edges $u v$ of $G$, i. e.,

$$
R(G)=\sum_{u v \in E(G)} \frac{1}{\sqrt{\operatorname{deg}(u) \operatorname{deg}(v)}}
$$

[^0]where $\operatorname{deg}(v)$ is the degree of a vertex $v$. Randić has shown that there exists a correlation of the Randić index with several physico-chemical properties of alkanes such as boiling points, chromatographic retention times, enthalpies of formation, parameters in the Antoine equation for vapor pressure, surface areas and others. More information about Randić index can be found in a survey [7] by Li and Shi or in a book [8] by Li and Gutman.

The Wiener index is a distance-based topological index defined as the sum of distances between all pairs of vertices in a graph and is denoted by $W(G)$. It was the first topological index in chemistry [13], introduced in 1945. Since then, the Wiener index has been used to explain various chemical and physical properties of molecules. However the Wiener index can be cubic in the number of vertices of a graph, hence for the sake of simplicity we use the average distance instead

$$
\operatorname{ad}(G)=W(G) /\binom{n}{2} .
$$

For a one vertex graph $K_{1}$ we assume $R\left(K_{1}\right)=\operatorname{ad}\left(K_{1}\right)=0$.
For the last two decades researchers are investigating extremal values and relations between topological indices. In 1988 Fajtlowicz [5] stated the following conjecture based on the computer program Graffiti.

Conjecture 1.1. For all connected graphs $G, R(G) \geq \operatorname{ad}(G)$.

Very recently Li and Shi [9] have proven Conjecture 1.1 when $\delta(G) \geq n / 5$ and $n \geq 15$. More recent results related to extremal values of the Randić index can be found in $[1,3,6,14,15]$.

In order to state the results of this paper let us introduce some notation. By $\operatorname{dist}_{G}(u, v)$ we denote the length of the shortest path between vertices $u$ and $v$ in a connected graph $G$. The length $\max _{u, v} \operatorname{dist}_{G}(u, v)$ is a diameter of a graph $G$, which we denote by $\operatorname{diam}(G)$. In particular for trees the diameter is the length of the longest simple path.

A star is a tree where at most one vertex is of degree greater than one. Note that a tree with at most three vertices is a star. Also notice that stars are precisely the trees with diameter at most two.

The main result of this paper is the following theorem.

Theorem 1.2. For any tree $T$ with $n$ vertices and $n_{1}$ leaves the following inequality holds:

$$
R(T) \geq \operatorname{ad}(T)+\max \left(0, \sqrt{n_{1}}-2\right) .
$$

Moreover if we consider the limit when $n$ goes to infinity, the inequality is sharp for stars.

As a consequence we prove Conjecture 1.1 for the case when $G$ is a tree.

## 2 Proof of Theorem 1.2

A vertex of degree one is called a pendent vertex or a leaf. By $\mathcal{T}_{n, n_{1}}$ we denote the set of all trees with exactly $n$ nodes and $n_{1}$ pendent vertices.

First we define a special class of trees, which will be frequently used in our paper. For $a, b \geq 1, n \geq a+b+2$ by $\mathrm{DC}(n, a, b)$ we denote a double comet, which is a tree composed of a path containing $n-a-b$ vertices with $a$ pendent vertices attached to one of the ends of the path and $b$ pendent vertices attached to the other end of the path. Thus, $\mathrm{DC}(n, a, b)$ has $n$ vertices and $a+b$ leaves, i. e., $\mathrm{DC}(n, a, b) \in \mathcal{T}_{n, a+b}$.


Figure 1: A double comet $\mathrm{DC}(n, a, b)$ for $n=16, a=5, b=4$.

We begin our proof with a lemma proving Theorem 1.2 for stars and paths.

Lemma 2.1. Let $T \in \mathcal{T}_{n, n_{1}}$ be a tree such that $n_{1} \leq 2$ (a path) or $n_{1}=n-1$ (a star). Then $R(T) \geq \operatorname{ad}(T)+\max \left(0, \sqrt{n_{1}}-2\right)$.

Proof. If $n \leq 2$ then the inequality trivially holds, hence we assume $n \geq 3$. If $T$ is a star then by direct calculations we obtain $\operatorname{ad}(T)=2-\frac{2}{n}$ and $R(T)=\sqrt{n_{1}}=\sqrt{n-1}$, whereas if $T$ is a path then $\operatorname{ad}(T)=\frac{n+1}{3}$ and $R(T)=\sqrt{2}+\frac{n-3}{2}$. In both cases the lemma holds.

Consequently from now on we assume that $T \in \mathcal{T}_{n, n_{1}}$ where $3 \leq n_{1} \leq n-2$, which means that we may use the following result proved by Liu et al. in Theorem 5 from [10].

Theorem 2.2. Let $T \in \mathcal{T}_{n, n_{1}}$, where $3 \leq n_{1} \leq n-2$. Then

$$
R(T) \geq \frac{n-n_{1}}{2}+\sqrt{n_{1}}+(1-\sqrt{2})\left(\frac{1}{\sqrt{2 n_{1}}}+\frac{1}{\sqrt{2}}\right) .
$$

Corollary 2.3. Using the assumption $n-2 \geq n_{1} \geq 3$ we obtain:

$$
R(T) \geq \frac{n-n_{1}}{2}+\sqrt{n_{1}}-0.462
$$

The flow of the proof is as follows. Firstly we show that each tree can be transformed into a double comet without decreasing the average distance and simultaneously preserving the number of vertices and leaves. Since in Corollary 2.3 we only use $n$ and $n_{1}$ disregarding the actual structure of a tree it is enough to show that for double comets we have $\frac{n-n_{1}}{2}+\sqrt{n_{1}}-0.462 \geq \operatorname{ad}(T)+\max \left(0, \sqrt{n_{1}}-2\right)$.

Lemma 2.4. Let $T \in \mathcal{T}_{n, n_{1}}$, where $3 \leq n_{1} \leq n-2$. There exists a double comet $T^{\prime}=\mathrm{DC}(n, a, b)$ for some $a, b \geq 1, a+b=n_{1}$ such that $\operatorname{ad}(T) \leq \operatorname{ad}\left(T^{\prime}\right)$.

Proof. Let us assume that contrary. Let $T \in \mathcal{T}_{n, n_{1}}$ be a non-double-commet tree such that for each double comet $T^{\prime}=\mathrm{DC}(n, a, b)$, where $a, b \geq 1, a+b=n_{1}$ we have $\operatorname{ad}\left(T^{\prime}\right)<\operatorname{ad}(T)$. Among all such trees from $\mathcal{T}_{n, n_{1}}$ we consider a tree with the largest diameter $\operatorname{diam}(T)$.

Let $d=\operatorname{diam}(T)$ and $D=v_{1} v_{2} \ldots v_{d}$ be a diameter of the tree $T$. Since $T$ is neither a star nor a double comet we have $d \geq 5$. Moreover, we know that $v_{1}, v_{d}$ are leaves and all neighbors of $v_{2}, v_{d-1}$ outside $D$ are pendents, because otherwise $D$ would not be a diameter. Furthermore, since $T$ is not a double comet there exists a vertex outside the diameter $x \notin D$ that is a neighbor of some vertex $v_{i}$ for $3 \leq i \leq d-2$. Observe that after a removal of the edge $x v_{i}$ the tree $T$ decomposes into exactly two parts. Let us denote by $T_{x}$ the part containing the vertex $x$ and by $T_{D}$ the other one. Similarly after removal of the edge $v_{i-1} v_{i}$ the tree $T_{D}$ decomposes into two parts $T_{D_{1}}, T_{D_{2}}$, where $v_{i-1} \in V\left(T_{D_{1}}\right)$ and $v_{i} \in V\left(T_{D_{2}}\right)$. Without loss of generality we may assume, that $\left|V\left(T_{D_{1}}\right)\right| \leq\left|V\left(T_{D_{2}}\right)\right|$, since otherwise we can reverse the diameter $D$.

Let $T^{\prime}$ be a tree $T$ with the edge $x v_{i}$ removed and $x v_{2}$ inserted (as in Fig. 2). We claim $\operatorname{ad}\left(T^{\prime}\right) \geq \operatorname{ad}(T)$. Observe that distances between vertices within $V\left(T_{D}\right)$ and within $V\left(T_{x}\right)$ in both trees $T, T^{\prime}$ are the same. Now we consider distances between vertices in $V\left(T_{D}\right)$ and $V\left(T_{x}\right)$. Let $a \in V\left(T_{D}\right), b \in V\left(T_{x}\right)$. Consider two cases:

- $a \in V\left(T_{D_{1}}\right)$, then $\operatorname{dist}_{T^{\prime}}(a, b) \geq \operatorname{dist}_{T}(a, b)-(i-2)$
- $a \in V\left(T_{D_{2}}\right)$, then $\operatorname{dist}_{T^{\prime}}(a, b) \geq \operatorname{dist}_{T}(a, b)+(i-2)$.

Since $\left|V\left(T_{D_{1}}\right)\right| \leq\left|V\left(T_{D_{2}}\right)\right|$, we have $W\left(T^{\prime}\right) \geq W(T)$ and so $\operatorname{ad}\left(T^{\prime}\right) \geq \operatorname{ad}(T)$. If $\left|V\left(T_{x}\right)\right|>1$ then $\operatorname{diam}\left(T^{\prime}\right)>\operatorname{diam}(T)$, hence the lemma holds for $T^{\prime}$ because $\operatorname{diam}\left(T^{\prime}\right)$ is larger. Which means that there exists a double comet $T^{\prime \prime}=D C(n, a, b)$, where $a+b=n_{1}$, such that $\operatorname{ad}\left(T^{\prime \prime}\right) \geq \operatorname{ad}\left(T^{\prime}\right) \geq \operatorname{ad}(T)$, a contradiction. Thus we know that $\left|V\left(T_{x}\right)\right|=1$. However we can continue in this manner moving pendent neighbors of vertices $v_{i}$ for $3 \leq i \leq d-2$ ending up with a double comet, thus completing the proof.

Obviously, the biggest value of $\operatorname{ad}(\operatorname{DC}(n, a, b))$, where $a+b=n_{1}$, is obtained when $a, b$ are as close as possible to $n_{1} / 2$, i. e., $a=\left\lfloor\frac{n_{1}}{2}\right\rfloor$ and $b=\left\lceil\frac{n_{1}}{2}\right\rceil$. These bounds are also observed in $[4,12]$ (see Theorem 31 in the survey [2]) but as the papers $[4,12]$ may not be easy to access and as our proof is short, for the sake of completness we decide to include it.


Figure 2: Ilustration of the transformation used in the proof of Lemma 2.4.

Lemma 2.5. Let $T \in \mathcal{T}_{n, k}$ be a double comet $\operatorname{DC}(n, a, b)$ for $a, b \geq 1$, where $3 \leq$ $a+b=n_{1} \leq n-2$. Then

$$
\frac{n-n_{1}}{2}+\sqrt{n_{1}}-0.462 \geq \operatorname{ad}(T)+\max \left(0, \sqrt{n_{1}}-2\right)
$$

Proof. Firstly we calculate the average distance for the double comet $T$ :

$$
\begin{aligned}
\binom{n}{2} \operatorname{ad}(T) & =W(T) \\
& =a b\left(n-n_{1}+1\right)+2\binom{a}{2}+2\binom{b}{2} \\
& +(a+b)\left(n-n_{1}\right)\left(n-n_{1}+1\right) / 2+\binom{n-n_{1}+1}{3}
\end{aligned}
$$

where $a b\left(n-n_{1}+1\right)+2\binom{a}{2}+2\binom{b}{2}$ comes from distances between leaves, $(a+b)(n-$ $\left.n_{1}\right)\left(n-n_{1}+1\right) / 2$ is the sum of distances between leaves and inner vertices and $\binom{n-n_{1}+1}{3}$ is the sum of distances between inner vertices. Using following inequalities:

$$
\begin{aligned}
2 a b+2\binom{a}{2}+2\binom{b}{2} & \leq(a+b)^{2}=n_{1}^{2} \\
a b & \leq \frac{n_{1}^{2}}{4} \\
\binom{n-n_{1}+1}{3} & \leq \frac{\left(n-n_{1}\right)^{3}}{6}
\end{aligned}
$$

we obtain

$$
\begin{aligned}
\operatorname{ad}(T)\binom{n}{2} & \leq \frac{n_{1}^{2}\left(n-n_{1}-1\right)}{4}+n_{1}^{2}+\frac{n_{1}\left(n-n_{1}\right)\left(n-n_{1}+1\right)}{2}+\frac{\left(n-n_{1}\right)^{3}}{6} \\
& =\frac{n_{1}^{2}\left(n-n_{1}\right)}{4}+\frac{3 n_{1}^{2}}{4}+\frac{n_{1}\left(n-n_{1}\right)^{2}}{2}+\frac{n_{1}\left(n-n_{1}\right)}{2}+\frac{\left(n-n_{1}\right)^{3}}{6} .
\end{aligned}
$$

For the sake of simplicity we put $x=n_{1}, y=n-n_{1}$.

$$
\operatorname{ad}(T)\binom{n}{2} \leq \frac{x^{2} y}{4}+\frac{3 x^{2}}{4}+\frac{x y^{2}}{2}+\frac{x y}{2}+\frac{y^{3}}{6}
$$

In order to prove the lemma it is enough to show the following inequality:

$$
\binom{n}{2}\left(\frac{n-n_{1}}{2}+\sqrt{n_{1}}-0.462-\max \left(0, \sqrt{n_{1}}-2\right)-\operatorname{ad}(T)\right) \geq 0 .
$$

We multiply the expression by 4 , put $x, y$ instead of $n_{1}, n-n_{1}$ and use the previously obtained inequality for $\operatorname{ad}(T)\binom{n}{2}$.

$$
\begin{array}{r}
4\binom{n}{2}\left(\frac{n-n_{1}}{2}+\sqrt{n_{1}}-0.462-\max \left(0, \sqrt{n_{1}}-2\right)-\operatorname{ad}(T)\right)= \\
2(x+y)(x+y-1)\left(\frac{y}{2}+\min (\sqrt{x}-0.462,1.538)\right)-4 \operatorname{ad}(T)\binom{n}{2} \geq \\
2\left(x^{2}+y^{2}+2 x y-x-y\right) \min (\sqrt{x}-0.462,1.538)+y^{3} / 3-3 x y-y^{2}-3 x^{2} .
\end{array}
$$

Now we consider two cases. Either $n_{1}=x=3$ or $n_{1}=x \geq 4$. When $x=3$ from the above expression we obtain:

$$
\begin{array}{r}
2\left(9+y^{2}+6 y-3-y\right) \min (\sqrt{3}-0.462,1.538)+y^{3} / 3-9 y-y^{2}-27 \\
2.5\left(y^{2}+5 y+6\right)+y^{3} / 3-y^{2}-9 y-27= \\
y^{3} / 3+1.5 y^{2}+3.5 y-12 .
\end{array}
$$

Using the assumption $n_{1} \leq n-2$ which is equivalent to $y \geq 2$ we obtain $y^{3} / 3+$ $1.5 y^{2}+3.5 y-12 \geq 11 / 3 \geq 0$.

Now we assume that $x \geq 4$, hence $\min (\sqrt{x}-0.462,1.538) \geq 1.5$. Since $x, y \geq 2$ then $x y \geq x+y$, and so we have:
$2\left(x^{2}+y^{2}+2 x y-x-y\right) 1.5+y^{3} / 3-3 x y-y^{2}-3 x^{2}=y^{3} / 3+2 y^{2}+3 x y-3(x+y) \geq 0$ what ends the proof of the lemma.

In Lemma 2.1 we prove Theorem 1.2 for paths and stars whereas by Lemmas 2.4 and 2.5 together with Corollary 2.3 we obtain the inequality for all other trees, thus completing the proof of Theorem 1.2.

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