

CONTEXT-FREE GRAMMAR AND DETERMINISTIC AUTOMATON APPROACHES
FOR SEQUENCE GENERATION IN COPOLYMERS. I. BINARY COPOLYMERS
WITH ULTIMATE EFFECT

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1. Abstract.

Linguistic pattern recognition techniques have been used for generating and identifying the near sequence in binary copolymers on the basis of the ultimate effect.

2. Introductory Notions in Linguistic Pattern Recognition.

Linguistic recognition uses the analogy between letters and the words within a phrase and the structure of a pattern. In this analogy any pattern (geometrical feature, digital print, image, etc.) [1, 2] may be decomposed into a number of component parts (called primitives) that have for correspondent the letters and the words of a language. On the other side the rules of a pattern composition may be related to the grammar rules of building a phrase. Therefore we come to the conclusion that defining a language and a grammar that describes a certain pattern class is possible. In order to exemplify let us consider the geometrical pattern of a square. The square can be decomposed (Figure 1) into a number of four

primitives a, b, c, d that define unit length segments. Therefore any sequence of four letters of abcd, bada, cdab, dabc type will identify the

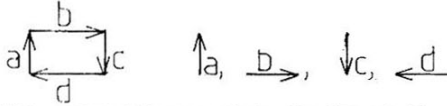


FIGURE 1. Primitives used in the description of a square.

square. Further on, the language L_1 that generates the set of the squares having sides of length equal to n units will be defined by the set:

$$L_1 = \left\{ a^n b^n c^n d^n \mid n \geq 1 \right\}$$

In order to specify the rules of building squares using elements of this language we use the theory of pattern languages that defines a grammar of the following pattern:

$$G = (V_N, V_T, P, S)$$

where:

- V_N - the set of the non-terminal symbols;
- V_T - the set of the terminal symbols;
- P - grammar rules;
- S - initial symbol.

The grammar elements that generate the patterns described by means of the L_1 language are:

$$V_T = \left\{ a, b, c, d \right\}$$

$$V_N = \left\{ S, A, B, C, D \right\}$$

$$P : S \rightarrow a A, \quad A \rightarrow a A, \quad B \rightarrow b B, \quad D \rightarrow d D$$

$$A \rightarrow B \quad B \rightarrow c C \quad D \rightarrow d$$

$$A \rightarrow b B \quad C \rightarrow d D$$

The operation through which we build the pattern, following the grammar rules specified by P is called derivation and is represented by the symbol " \Rightarrow ". In the case of a square having unit sides, the derivation may have the pattern:

$$S \Rightarrow a A \Rightarrow a b B \Rightarrow a b c C \Rightarrow a b c D \Rightarrow a b c d$$

Such a grammar is known under the name of context-free grammar. If we wish to represent a quadrilateral, whose primitives are described in Figure 2, the language will have the pattern:

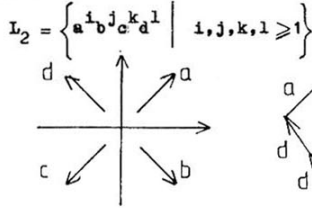


FIGURE 2. Primitives used in the description of a quadrilateral.

In this case the grammar is stochastic and has the pattern:

$$G_S = (V_N, V_T, P_S, S)$$

where P_S is a finite set of stochastic rules.

For a context-free stochastic grammar a rule in P_S has the pattern:

$$A_i \xrightarrow{P_{ij}} \alpha_j, A_i \in V_N, \alpha_j \in (V_N \cup V_T)$$

where P_{ij} represents the production probability.

In the case of a quadrilateral the grammar rules are given by:

$$P_S: S \rightarrow aA, A \xrightarrow{P_{11}} aA, B \xrightarrow{P_{22}} bB, C \xrightarrow{P_{33}} cC, D \xrightarrow{P_{44}} dD$$

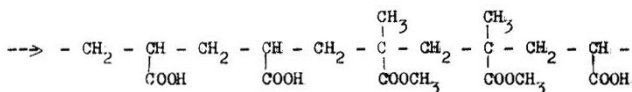
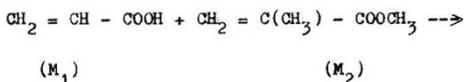
$$A \xrightarrow{P_{12}} B, B \xrightarrow{P_{23}} C, C \xrightarrow{P_{34}} D, D \rightarrow d$$

Supposing the probability distribution is uniform within the interval (0, 1), the derivation rules for a quadrilateral are:

$$S \Rightarrow aA \xRightarrow{P_{12}} aB \xRightarrow{P_{22}} abB \xRightarrow{P_{23}} abcC \xRightarrow{P_{33}} abccC \xRightarrow{P_{34}} abccD \xRightarrow{P_{44}} abccdd$$

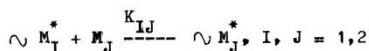
3. Generating Copolymer Sequences.

A copolymer represents a chemical macromolecular compound having an aleatory repetitive structure in which the setting-up of mer sequences (the mer being a structural elementary unit of a macromolecular chain) as well as their nature and number depend on the conditions in which the chemical reaction of copolymer formation takes place (temperature, pressure, solvent catalyst, etc.). The nature and the reactivity of the monomers also influence the mer sequence. As there is a close link between the copolymer structure and its properties, any theoretical approach in determining the structure on the ground of the monomer reactivity may prove to be extremely useful as it might allow the synthesis of macromolecular compounds having an a-priori structure. Therefore, determining sequence distribution is important for understanding physical properties of the copolymer, for determining reactivity parameters of monomers and for discriminating between many possible reaction mechanisms. In order to exemplify we present a binary copolymerization reaction between acrylic acid (M_1) and methyl methacrylate (M_2) monomers, as well as a possible $M_1M_1M_2M_2M_1$ sequence in the copolymer chain:



Several models have been developed to explain the way in which copolymerization reactions take place and to allow copolymer structure determination according to monomer reactivity [4]. All these studies take into account the appearance probability distribution of a mer conditioned by the growing macro-radical structure (the ultimate mer-model

and the ultimate effect, the penultimate mer-models and penultimate effect etc.) as well as the conditions of setting up a copolymerization reaction, expressed through r_1 and r_2 copolymerization reactivity ratios. A binary copolymerization reaction (two monomers) with an ultimate effect (that is only terminal monomer units - mers - alter the rate values) imply four elementary growing reactions [3]:



The ratios $r_I = K_{II} / K_{IJ}$, $I, J = 1, 2, I \neq J$ are called reactivity ratios. Supposing that:

i) rate values do not depend on the length of the growing macroradical;

ii) we can neglect initiation or ending processes (that is we consider that the polymer length is so large that it does not influence initiation or ending macroradical growing), the probabilities of transition P_{II} from the state M_I^* to the state M_J^* can easily be calculated thus:

$$P_{II} = \frac{K_{II} (M_I^*) (M_I)}{K_{II} (M_I^*) (M_I) + K_{IJ} (M_I^*) (M_J)} =$$

$$= \frac{K_{II} (M_I)}{K_{II} (M_I) + K_{IJ} (M_J)} = \frac{r_I (M_I)}{r_I (M_I) + (M_J)} = \frac{r_I}{r_I + C}$$

where () designs the initial monomer concentration and $C = (M_J) / (M_I)$. P_{IJ} can be obtained with the following relation:

$$P_{IJ} = 1 - P_{II} = \frac{r_I}{r_I + C}$$

The authors intend to develop a linguistic copolymer sequence generation and recognition system. They have thus proposed the interpre-

tation of a copolymer sequence as sentences of different lengths in which the mers are letters that group in words formed of one or many mers of the same type. Here are the steps we have made:

1° Choosing the primitives is imposed by the mer type that composes the copolymer. As in the case of binary copolymerization the monomer units are of two different types, that have been chosen as primitives, the mers (the primitives) being noted with M_1 and M_2 . As noted above, in the case of the copolymer acrylic acid / methyl methacrylate, primitive M_1 will be represented by the mer - $\text{CH}_2 - \text{C}(\text{CH}_3) (\text{COOCH}_3) -$.

2° Copolymer structure representation is made under the form of a primitive string. Thus, for the situation mentioned above a sequence in a primitive string is $M_1 M_1 M_2 M_1 M_2 M_2 \dots$

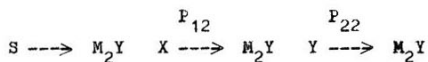
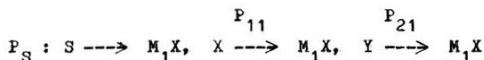
3° A context-free stochastic grammar of the pattern below has been chosen for describing and generating sentences that represent such structures:

$$G_S = (V_N, V_T, P_S, S)$$

where:

$$V_N = \left\{ S, X, Y \right\}$$

$$V_T = \left\{ M_1, M_2 \right\}$$



$$P_{11} = \frac{r_1}{r_1 + C}, \quad P_{12} = 1 - P_{11}, \quad P_{22} = \frac{r_2}{r_2 + C}$$

$P_{21} = 1 - P_{22}$; r_1 and r_2 are copolymerization rate values.

The grammar thus defined has been implemented in a program for

generating copolymer sequence in which the copolymerization degree represents the sequence building process stopping criterion [4 - 11]. The probability of generating the string that represents the copolymer appears as the product of all probability productions associated with the production used in mer generating. We have considered that if the last mer is of the M_1 or M_2 type the appearance probability of another M_1 or M_2 mer to be attached to the growing chain is uniformly distributed; the appearance probabilities are P_{11} , P_{12} respectively if the last mer is of the M_2 type. A sequence generated according to this grammar is shown in Figure 3, and the GENERATE-STRING sequence generation program is presented in the Appendix 1.

$M_1-M_2-M_1-M_1-M_2-M_1-M_2-M_2-M_1-M_2-M_2-M_1-M_1-M_2-M_1-M_2-M_2-M_2-M_1-M_2-M_1-M_1-M_1-M_1-M_1-$
 $M_1-M_1-M_2-M_1-M_1-M_1-M_2-M_1-M_2-M_1-M_1-M_1-M_2-M_2-M_2-M_2-M_2-M_2-M_2-M_2-M_1-M_1-M_1-M_2-$
 $M_1-M_1-M_2-M_1-M_2-M_2-M_1-M_2-M_2-M_2-M_1-M_2-M_1-M_2-M_1-M_1-M_1-M_1-M_2-M_1-M_1-M_1-M_1-$
 $M_1-M_2-M_2-M_1-M_1-M_2-M_1-M_2-M_1-M_1-M_1-M_1-M_1-M_1-M_1-M_1-M_2-M_1-M_1-M_1-M_1-M_1-$
 $M_1-M_1-M_1-M_1-M_1-M_1-M_1-M_1-M_1-M_2-M_2-M_1-M_1-M_1-M_1-M_1-M_1-M_1-M_1-M_1-M_2-M_1-M_1-$
 $M_1-M_2-M_1-M_2-M_2-M_2-M_2-M_1-M_1-M_2-M_1-M_1-M_1-M_2-M_2-M_1-M_1-M_1-M_1-M_1-M_2-M_1-M_1-$
.....
.....
 $M_2-M_1-M_2-M_1-M_2-M_1-M_1-M_2-M_1-M_1-M_1-M_1-M_1-M_1-M_1-M_1-M_1-M_2-M_2-M_1-M_2-M_1-M_2-M_1-$
 $M_2-M_1-M_1-M_1-M_1-M_1-M_1-M_2-M_1-M_1-M_2-M_1-M_1-M_2-M_1-M_2-M_2-M_1-M_1-M_2-M_1-M_2-M_1-$

FIGURE 3. ($M_1 * M_2$) 2-vinyl-pyridine*styrene copolymer partial string sequence generated according to G_3 .

4. Principles of Copolymer Linguistic Recognition.

If a pattern class can be described through a determinist language then the procedure of recognizing a pattern to belong to a certain class is made by a determinist automaten. This accepts as the input the

string describing the pattern, and verifies the belonging of a pattern to a certain class with the help of grammar rules used in building a pattern. For the pattern generated by the L language $L = \left\{ a^n b^n c^n d^n \mid n \geq 1 \right\}$ the determinist automaton has the pattern:

$$A = (\Sigma, Q, \rightarrow, q_0, F)$$

where:

$\Sigma = \{a, b, c, d\}$ and is the alphabet;

$Q = \{q_0, q_1, q_2, q_3\}$ the state set;

$F = \{q_3\}$ final state set;

q_0 is the initial state.

δ is an application on the $Q \times F$ set with values in the Q set.

The transition diagram has the pattern:

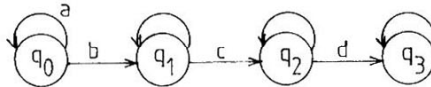


FIGURE 4. The transition diagram of the A automaton.

For the A automaton, δ will have the values:

$$\delta(q_0, a) = \{q_0, q_1\}$$

$$\delta(q_1, b) = \{q_1, q_2\}$$

$$\delta(q_2, c) = \{q_2, q_3\}$$

$$\delta(q_3, d) = \{q_3\}$$

In the case of pattern recognition whose generation was made with the help of a stochastic grammar, the most simple procedure of recognition can be made with the help of a non-determinist automaton. This solution has been adopted by the authors for the recognition (identification) of a copolymer. The next non-determinist automaton has been defined:

$$C = (\Sigma, Q, q_0, F)$$

where:

$$\Sigma = \{M_1, M_2\}$$

$$Q = \{q_0, q_1, q_2\}$$

$$F = \emptyset$$

$$\delta(q_0, M_1) = \{q_1, q_2\}$$

$$\delta(q_1, M_1) = \{q_1\}$$

$$\delta(q_2, M_2) = \{q_2\}$$

$$\delta(q_2, M_1) = \{q_1\}$$

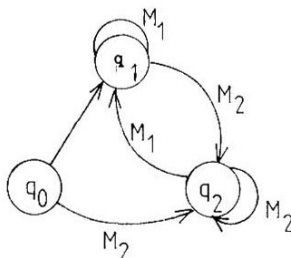


FIGURE 5. The transition diagram of the C automaton.

The transition diagram of the C automaten is given in Figure 5. The automaten was implemented through a computer program. This program accepts as input the string that describes the copolymer, and affords the rate values as output. We have calculated the rate values r_1 and r_2 of the copolymer that results from the methyl methacrylate / chloroprene copolymerization:

$$r_1 = 0.073 \quad \text{and} \quad r_2 = 5.852$$

(the experimental [12] ones having the following values: $r_1 = 0.08$ and $r_2 = 5.1$). The RECOGNITION-STRING program is presented in the Appendix 2.

5. Results.

In order to exemplify (See Table 1) we offer below the results obtained in the case of copolymerization of acrylic acid (M_1) with methyl methacrylate (M_2).

	M_1		M_1 in the obtained copolymer	
	initial	experimental	Markov (M)	Automaton (C)
1.	88.0	91.1	90.8	92.0
2.	88.1	85.7	84.9	85.4
3.	72.4	77.2	77.2	76.7
4.	71.6	75.6	76.4	75.5
5.	58.9	61.5	63.9	61.2
6.	58.7	62.2	63.7	61.1
7.	51.0	56.1	55.6	53.4
8.	45.8	49.1	50.0	46.6
9.	34.0	41.1	36.9	33.0
10.	30.5	32.6	33.0	28.7
11.	20.2	25.1	21.5	20.2

TABLE 1.

The agreement between the compositions calculated using a Markov chain model of the first order and the stochastic linguistic grammar described above is very good, as it is shown in the following linear regression equations:

$$\% M_{1, \text{exp}} = 7.658 + 0.904 \% M_{1, C}$$

$$(r = 0.998, s = 1.371, F = 981.779)$$

$$\% M_{1, \text{exp}} = 2.854 + 0.957 \% M_{1, M}$$

$$(r = 0.99, s = 1.759, F = 594.962)$$

where r represents the correlation coefficient, s the standard deviation and F the Fisher statistics.

Acknowledgement.

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APENDIX 1

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C *****
C *****
C *****
C *****          G E N E R A T E - S T R I N G          *****
C *****
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C *****
```

C GENERATE-STRING IS NEW VERSION OF MEMORY-3

BLOCK DATA

COMMON/TMER/PKTA,PKTB,ELK

DATA PKTA,PKTB,ELK/2HA-,2HB-,1H /

END

INTEGER SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NEA, NEB, RA, RB

COMMON N

COMMON/KONST/SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NEA, NEB, RA, RB, I

COMMON/POLY/POLY(5000)/DIVID/ALPHA, BETA, C, R1, R2

COMMON/TABEL/GTAB(5,16),J/RANDOM/IA, IB, X

DIMENSION NAME(40)

L=108

1 READ(105,900,END=8)(NAME(K),K=1,40)

READ(105,910)MER, N, R1, R2, NA, NB, FRMOL

WRITE(L,940)

WRITE(L,900)(NAME(K),K=1,40)

WRITE(L,950)N, R1, R2, FRMOL, NA, NB

CALL INIT

CALL ERROR(MC)

IF(MC.EQ.1)GO TO 1

IF(MER.EQ.1)CALL INITA

IF(MER.EQ.2)CALL INITB

2 CALL FARI

```
3    CALL ERRORP(MC)
      IF(MC.EQ.1)GO TO 1
      IF(MC.EQ.1)GO TO 1
4    CALL MAS
      CALL ALEAT
      CALL NMAS
      IA=IB
      CALL ERRORR(MC)
      IF(MC.EQ.1)GO TO 1
      I=I+1
      IF(I.GT.N)CALL OUT(MC)
      IF(MC.EQ.1)GO TO 1
      GO TO (5,6),SW
5    CALL GENER1
      GO TO 7
6    CALL GENER2
7    CALL REEW(MC)
      IF(MC)2,3,4
8    STOP
900  FORMAT(40A2)
910  FORMAT(I2,I6,2F14.8,3I6)
940  FORMAT(1H0///)
950  FORMAT(5X,'POLYMERISATION DEGREE=' ,I6/
1     SX,'          R1=' ,F14.8/
2     SX,'          R2=' ,F14.8/
3     SX,'EVALUATION MODE =' ,I6/
4     SX,'          NA=' ,I6/
4     SX,'          NB=' ,I6//)
      END)
      SUBROUTINE REEW(MC)
```

```
INTEGER SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB
COMMON/KONST/SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB, I
COMMON/POLYM/POLYM(5000)/DIVID/ALPHA, BETA, C, R1, R2
COMMON/TMER/PKTA, PKTB, BLK
IF(((ALPHA.EQ.0.).AND.(BETA.EQ.0.)).OR.
1 ((ALPHA.EQ.1.).AND.(BETA.EQ.1.))) GO TO 1
IF((NTAS+NTBS).LT.FRMO) GO TO 1
NA=NA-NTAS
NB=NB-NTBS
NTAS=0
NTBS=0
IF((NB.LT.0).OR.(NA.LT.0)) GO TO 2
IF(NA.EQ.0) GO TO 3
IF(NB.EQ.0) GO TO 4
MC=-1
RETURN
2 MC=0
RETURN
3 CALL ENDA
1 MC=1
RETURN
4 CALL ENDB
MC=1
RETURN
END
SUBROUTINE JN11
INTEGER SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB
COMMON/KONST/SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB, I
COMMON/POLYM/POLYM(5000)/DIVID/ALPHA, BETA, C, R1, R2
COMMON/TABEL/UTAB(5,16),J/RANDOM/IA,IB,X
```

```
COMMON/TMER/PKTA,PKTB,BLK
I=1
J=1
IA=65539
NRA=0
NRB=0
NBA=0
NBB=0
NTAS=0
NTBS=0
DO 1 L=1,5
DO 1 K=1,16
1  GTAB(L,K)=BLK
DO 2 L=1,20
CALL MAS
CALL ALEAT
CALL NMAS
2  IA=IB
RETURN
ENTRY INITA
POLYM(1)=PKTA
SW=1
NRA=1
NTAS=1
RETURN
ENTRY INITB
POLYM(1)=PKTB
SW=2
NRB=1
NTBS=1
```

RE TURN

END

SUBROUTINE PART

INTEGER SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB

COMMON/KONST/SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB, I

COMMON/POLYM/POLYM(5000)/DIVID/ALPHA, BETA, C, R1, R2

C=FLOAT(NA)/FLOAT(NB)

ALPHA=R1/(1./C+R1)

BETA=1.-R2/(C+R2)

RETURN

END

SUBROUTINE ERROR(MC)

INTEGER SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB

COMMON N

COMMON/KONST/SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB, I

COMMON/POLYM/POLYM(5000)/DIVID/ALPHA, BETA, C, R1, R2/RANDOM/ IA, IB, X

MC=0

L=108

IF(N.LE.0) GO TO 1

IF(NB.LE.0.OR.NA.LE.0) GO TO 2

IF(FRMOL.LE.0) GO TO 3

IF(MER.LT.1.OR.MER.GT.2) GO TO 4

RETURN

1 WRITE(L,900)N

MC=1

RETURN

2 WRITE(L,910)NA,NB

MC=1

RETURN

3 WRITE(L,920)FRMOL


```
      MC=1
      RETURN
4    WRITE(L,930)MER
      MC=1
      RETURN
900  FORMAT(///5X,'*ERROR* N=',I6)
910  FORMAT(///5X,'*ERROR* NA=',I6,' NB=',I6)
920  FORMAT(///5X,'*ERROR* EVALUATION MODE=',I6)
930  FORMAT(///5X,'*ERROR* MER=',I2)
      ENTRY ERRORP(MC)
      L=108
      MC=0
      IF((ALPHA.LT.0.).OR.(ALPHA.GT.1.).OR.
1    1(BETA.LT.0.).OR.(BETA.GT.1)) GO TO 5
      RETURN
5    WRITE(L,940)ALPHA,BETA
      MC=1
      RETURN
940  FORMAT(///5X,'*ERROR* ALPHA=',F14.8,' BETA=',F14.8)
      ENTRY ERRORR(MC)
      MC=0
      L=108
      IF(X.LE.0.OR.X.GE.1.) GO TO 6
      RETURN
6    WRITE(L,960)X
      MC=1
      RETURN
960  FORMAT(///5X,'*ERROR* GENERATOR DEFECT X=',F14.7)
      END
      SUBROUTINE GENER1
```

```
INTEGER SW,FORMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB
COMMON/KONST/SW, FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, EBB, RA, RB, I
COMMON/POLYM/POLYM(5000)/DIVID/ALPHA, BETA, C, R1, R2
COMMON/TABEL/GTAB(S, 16), J/RANDOM/IA, IB, X
COMMON/TMER/PKTA, PKTB, BLK
1  IF(NA.GT.NTAS) GO TO 2
   IF(NA.EQ.0.AND.NTAS.EQ.0) GO TO 2
   CALL ENDA
   GO TO 3
2  IF(X.LE.ALPHA) GO TO 5
3  IF(NRA.EQ.1.OR.NRA.EQ.0) GO TO 4
   RA=NA-NTAS
   RB=NB-NTBS
   NBA=NBA+1
   GTAB(1, J)=PKTA
   GTAB(2, J)=NRA
   GTAB(3, J)=NBA
   GTAB(4, J)=RA
   GTAB(5, J)=RB
   J=J+1
   IF(J.GT.14) CALL GRTAB
4  POLYM(I)=PKTB
   SW=2
   NRA=0
   NRB=NRB+1
   NTBS=NTBS+1
   RETURN
5  POLYM(I)=PKTA
   NRA=NRA+1
   NTAS=NTAS+1
```

```
RETURN
END
SUBROUTINE GENER2
INTEGER SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB
COMMON/KONST/SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB, I
COMMON/POLYM/POLYM(5000)/DIVID/ALPHA, BETA, C, R1, R2
COMMON/TABEL/GTAB(5,16), J/RANDOM/IA, IB, X
COMMON/TMER/PKTA, PKTB, BLK
1  IF(NB.GT.NTBS) GO TO 2
   IF(NB.EQ.0.AND.NTBS.EQ.0) GO TO 2
   CALL ENDB
   GO TO 3
2  IF(X.LE.BETA) GO TO 3
   POLYM(1)=PKTB
   NRB=NRB+1
   NTBS=NTBS+1
   RETURN
3  IF(NRB.EQ.1.OR.NRB.EQ.0) GO TO 4
   RA=NA-NTAS
   RB=NB-NTBS
   NBB=NBB+1
   GTAB(1,J)=PKTB
   GTAB(2,J)=NRB
   GTAB(3,J)=NBB
   GTAB(4,J)=RA
   GTAB(5,J)=RB
   J=J+1
   IF(J.GT.14) CALL GGTAB
4  POLYM(I)=PKTA
   SW=1
```

```
NRB=0
NRA=NRA+1
NTAS=NTAS+1
RETURN
END
SUBROUTINE  OUT(MC)
INTEGER  SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB
INTEGER  G(2,40), A, B
COMMON  N
COMMON/KONST/SW,FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NBA, NBB, RA, RB
COMMON/POLYM/POLYM(5000)/DIVID/ALPHA, BETA, C, R1, R2
COMMON/TABEL/GTAB(5,16), J/RANDOM/IA, IB, X
COMMON/TMER/PKTA, PKTB, BLK
MC=0
L=108
IF (NRA.EQ.1.OR.NRA.EQ.0) GO TO 1
RA=NA-NTAS
RB=NB-NTBS
NBA=NBA+1
GTAB(1,J)=PKTA
GTAB(2,J)=NRA
GTAB(3,J)=NBA
GTAB(4,J)=RA
GTAB(5,J)=RB
1  IF(NRB.EQ.1.OR.NRB.EQ.0) GO TO 2
RA=NA-NTAS
RB=NB-NTBS
NRB=NRB+1
GTAB(1,J)=PKTB
GTAB(2,J)=NRB
```

```
GTAB(3,J)=NBB
GTAB(4,J)=RA
GTAB(5,J)=RB
2  IF(J.GT.1) CALL GRTAB
   IF(J.EQ.1.AND.(GTAB(1,J).EQ.PKTA.OR.GTAB(1,J).EQ.PKTB))CALL GRTAB
   DO 3 I=1,2
   DO 3 J=1,40
3  G(I,J)=0
   NJ=0
   A=0
   B=0
   IC=-1
   IF(POLYM(1).EQ.PKTA)IC=1
   DO 20 I=1,N
   IF(POLYM(1).EQ.PKTA) GO TO 4
   GO TO 8
4  IF(IC.EQ.-1)GO TO 5
   A=A+1
   GO TO 20
5  IF(B.GT.40) GO TO 7
   G(2,B)=G(2,B)+1
6  IC=1
   B=0
   GO TO 4
7  G(2,40)=G(2,40)+1
   GO TO 6
8  IF(IC.EQ.1) GO TO 9
   B=B+1
   GO TO 20
9  IF(A.GT.40) GO TO 11
```

```
G(1,A)=G(1,A)+1
10  IC=-1
    A=0
    NJ=NJ+1
    GO TO 8
11  G(1,40)=G(1,40)+1
    GO TO 10
20  CONTINUE
    IF(IC.EQ.-1) GO TO 30
    IF(A.GT.40) A=40
    G(1,A)=G(1,A)+1
    GO TO 40
30  IF(B.GT.40) B=40
    G(1,B)=G(1,B)+1
40  CONTINUE
    WRITE(L,920)(J,J=1,40),((G(I,J),J=1,40),I=1,2)
920  FORMAT(/5X,'CONFIGURATION'/5X,12(' ')/
14X,'N I',40I3/4X,123(' ')/4X,'A I',40I3/4X,'B I',40I3)
C   WRITE(6,900)(POLYM(I),I=1,N)
    WRITE(L,910)NBA,NBB,NJ
900  FORMAT(/5X,'THE MOST PROBABLE SHAPE OF THE MOLECULE'
1     /5X,'-----'
2     //(6X,30A2))
910  FORMAT(/5X,'NUMBER OF GROUPS A=',I7
1     /5X,'NUMBER OF GROUPS B=',I7
1/5X,'NUMBER OF GROUPS A-B =',I7)
    MC=1
    RETURN
    END
SUBROUTINE GR1AB
```

```
COMMON/TABEL/GTAB(5, 16), J/RANDOM/IA, IB, X
COMMON/POLYM/POLYM(5000)/DIVID/ALPHA, BETA, C, R1, R2
COMMON/TMER/PKTA, PKTB, BLK
L=108
C
WRITE(6, 900) ((GTAB(II, JJ), JJ=1, 14), II=1, 5)
DO 1 II=1, 5
DO 1 JJ=1, 14
1
GTAB(II, JJ)=BLK
J=1
RETURN
900
FORMAT(/2X, 130(---)/2X, /! TYPE          !',
1
      14(---, A2, ---), --- !',
2
      /2X, 130(---)/2X, /! LENGTHS          !',
3
      14(---, F5.0), --- !',
4
      /2X, 130(---)/2X, /! POSITION          !',
5
      14(---, F5.0), --- !',
6
      /2X, 130(---)/2X, /! NUMBER OF MERS A IN FGED !',
7
      14(---, F5.0), --- !',
8
      /2X, 130(---)/2X, /! NUMBER OF MERS B IN FELD !',
9
      14(---, F5.0), --- !',
9
      /2X, 130(---))
END
SUBROUTINE ENDB
INTEGER SW, FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NPA, NPB, RA, RB
COMMON/KONST/SW, FRMOL, MER, NA, NB, NTAS, NTBS, NRA, NRB, NPA, NPB, RA, RB, I
COMMON/TABEL /GTAB(5, 16), J/RANDOM/IA, IB, X
COMMON/POLYM/POLYM(5000)/DIVID/ALPHA, BETA, C, R1, R2
L=108
ALPHA=1.0
BETA=1.0
```

```
NTBS=0
FRMOL=24*N
NB=0
WRITE(L,900)
900  FORMAT(5X,'*** NUMBER OF MERK B = 0 **')
RETURN
END
SUBROUTINE ENDA
INTEGER SW,FRMOL,MER,NA,NB,NTAS,NTBS,NRA,NRB,NBA,NBB,RA,RB
COMMON/TABEL/GTAB(5,16),J/RANDOM/IA,IB,X
COMMON/KONST/SW,FRMOL,MER,NA,NB,NTAS,NTBS,NRA,NRB,NBA,NBB,RA,RB,I
COMMON/POLYM/POLYM(5000)/DIVID/ALPHA,BETA,C,R1,R2
L=108
ALPHA=0.0
BETA=0.0
NTAS=0
FRMOL=24*N
NA=0
WRITE(L,900)
900  FORMAT(5X,'*** NUMBER OF MERK A = 0 **')
RETURN
END
SUBROUTINE ALEAT
COMMON/RANDOM/IA,IB,X
CALL RAN
RETURN
END
SUBROUTINE RAN
COMMON/RANDOM/IA,IB,X
II=IA+65537
```



```
IF (IB.GE.0) GO TO 1
IB=IB+2147483647+1
1 X=IB
X=X*0.456613E-09
IA=IB
RETURN
END

      ENTSG MAS,NMAS
      COMPILE ASSIRIS
MASKEI   CSECT    P
          DEF      MAS,NMAS
MASQUE   DATA,4,4 X'50400000'
MAS      LDTM,13  MASQUE
          BRU      *32
NMAS     LDTM,13  MASQUE+1
          BRU      *32
END
```

APENDIX 2

```
C *****
C *****
C *****
C *****      IDENTIFICATION - STRING      *****
C *****
C *****
C *****
C *****

C  IDENTIFICATION-STRING IS MEMORY-8 PROGRAM

C          COMPUTATION OF THE REACTIVITY RATIOS IN BINARY

C          IRREVERSIBLE COPOLYMERISATION WITH ULTIMATE EFFECT.

C  INPUT DATA:

C          THE CARD CONTAINS:

C          NA - NUMBER OF MER A.

C          NB= - NUMBER OF MER B.

C          NJ - NUMBER OF GROUPS A-B.

      DIMENSION T(10000),R(10000)

      DATA A,B/1HA,1HB/

      CALL ASSIGN (4,'CR: ')

      READ(4,100) NA,NB,NJ

      CALL ASSIGN (6,'LP ')

      N=NA+NB

      CALL GEN(T,R,N)

      DO 3 KA=1,N,5

      DO 3 KB=1,N,5

      CALL GENM(T,R,N,KA,KB,NAB,NAI,NBI)

      IF((NA.GT.NAI-3.AND.NA.LT.NAI+3).AND.

1(NAB.GT.NJ-3.AND.NAB.LT.NJ+3)) GO TO 5

      GO TO 3

5 CALL PRINT(T,KA,KB,NAB,NAI,NBI)

3 CONTINUE

STOP
```

```
100 FORMAT(3I5)
      END
      SUBROUTINE PRINT(T,KA,KB,KAB,NAI,NBI)
      DIMENSION T(1)
      A=T(KA)
      B=T(KB)
      R1=A/(1.-A)
      R2=(1.-B)/B
      WRITE(6,100) A,B,R1,R2,KAB,NAI,NBI
100  FORMAT(SX,'A=',F10.8,SX,'B=',F10.8,SX,'R1=',F10.5,SX,'R2=',
      1F10.5/
      24X,'NJ=',I5,2X,'NAI=',I5,2X,'NBI=',I5)
      RETURN
      END
      SUBROUTINE GEN(T,R,N)
      DIMENSION T(1),R(1)
      DATA IA/65539/
      DO 1 I=1,100000
      CALL ALEAT(IA,IB,X)
1    IA=IB
      DO 2 I=1,N
      CALL ALEAT(IA,IB,X)
      IA=IB
      T(I)=X
2    R(I)=X
      CALL RANCRU(T,N,N)
      RETURN
      END
      SUBROUTINE GENFIT(R,N,KA,KB,NAB,NAI,NBI)
      DIMENSION T(1),R(1)
```

```
NAI=0
NBI=0
NAB=0
IAE=1
A=T(KA)
B=T(KB)
DO 5 J=1,N
GO TO (1,3), IAB
1 IF(R(I).LE.A) GO TO 2
NBI=NBI+1
NAB=NAB+1
IAB=2
GO TO 5
2 NAI=NAI+1
GO TO 5
3 IF(R(I).LE.B) GO TO 4
NBI=NBII+1
GO TO 5
4 NAI=NAI+1
IAB=1
5 CONTINUE
RETURN
END
SUBROUTINE RANCRU(Z,N,NP)
DIMENSION Z(NP)
NI=N-1
IF(Z(1).LE.Z(2)) GO TO 1
A=Z(2)
Z(2)=Z(1)
Z(1)=A
```

```
1 DO 4 I=2,N1
  IF(Z(I).LE.Z(I+1)) GO TO 4
  A=Z(I+1)
  Z(I+1)=Z(I)
  I1=I+1
  DO 2 J=1,I1
    IF(Z(I-J).LE.A) GO TO 3
    Z(I-J+1)=Z(I-J)
2 CONTINUE
  Z(1)=A
  GO TO 4
3 Z(I-J+1)=A
4 CONTINUE
  RETURN
  END
  SUBROUTINE ALEAT(IA,IB,Z)
  DATA I1,I2/0,0/
  Z=RAN(I1,I2)
  RETURN
  END
```