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The Minimum Forcing and Anti–Forcing Numbers of Convex Hexagonal Systems^{*} Yaxian Zhang, Heping Zhang[†]

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Abstract

A convex hexagonal system (CHS) is a hexagonal system whose inner dual has the convex polygonal boundary. The minimum forcing number of a graph G is the smallest cardinality of a matching of G contained in a unique perfect matching. The minimum anti-forcing number of G is the smallest cardinality of an edge subset of G whose deletion results in a graph with exactly one perfect matching. In this paper, we proved that for any convex hexagonal system $H(a_1, a_2, a_3)$ with a perfect matching, its minimum forcing and anti-forcing numbers are both equal to min $\{a_1, a_2, a_3\}$ by applying perfect path systems.

1 Introduction

A *hexagonal system*, also called *benzenoid system*, is a 2-connected plane bipartite graph so that every interior face is a regular hexagon. A *perfect matching* of a graph is a set of disjoint edges covering all vertices of it. This concept coincides with that of a Kekulé structure in organic chemistry. The carbon-skeleton of a benzenoid hydrocarbon may be represented by a hexagonal system with a perfect matching. Kekulé structures of hexagonal systems have been extensively investigated; for example, see [5, 7].

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For a perfect matching M in a graph G, an edge subset $S \subseteq M$ is called a *forcing set* of M in G if G - V(S) contains exactly one perfect matching. The smallest cardinality over all forcing sets of M is the *forcing number* of M, denoted by f(G, M). The minimum and maximum values of forcing numbers of all perfect matchings of G are called the minimum and maximum forcing number of G, respectively, denoted by f(G) and F(G). The concept of "forcing number" in perfect matchings was first introduced by Randić and Klein [9] under name "innate degree of freedom", and then renamed by Harary et al. [8].

Forcing number of a Kekulé structure of a benzenoid system is related closely to resonance theory. According to Clar's aromatic sextet theory [4], for any two isometric benzenoid hydrocarbons, the one with larger Clar number is more stable. Xu et al. [19] showed that the maximum forcing number of a hexagonal system H is equal to its Clar number (or resonant number). Further Zhou and Zhang [27] showed that for each perfect matching M of H with the maximum forcing number, the maximum set of disjoint M-alternating hexagons has the size equal to the Clar number. For polyomino graphs and (4,6)-fullerenes, the former assertion also holds (cf. [16, 28]). Zhang and Li [21] characterized a hexagonal system with a forcing edge. Wang et al. [18] gave a linear algorithm to compute the minimum forcing number of a toroidal polyhex. Recently, Diwan [6] proved that for a hypercube Q_n , $f(Q_n) = 2^{n-2}$, which was ever conjectured by Pachter and Kim [14]. Afshani et al. [2] proved that determining the minimum forcing number is NP-complete for bipartite graphs with maximum degree 4.

As early as 1997, Li [12] characterized hexagonal systems with a forcing single edge (i.e. anti-forcing edge). In 2007, Vukičević and Trinajstić [17] proposed the anti-forcing number of a graph. Lei et al. [11] and Klein and Rosenfeld [10] independently generalized the idea of "anti-forcing" to every perfect matching. For a perfect matching M in a graph G, a set S of edges of G not in M is called an *anti-forcing set* of M if G - S has a unique perfect matching M. The cardinality of a smallest anti-forcing set of M is called the anti-forcing number of M, denoted by af(G, M). The minimum (resp. maximum) anti-forcing number of G is denoted by af(G) (resp. Af(G)). Lei et al. [11] and Shi et al. [16] respectively proved that the maximum anti-forcing numbers of hexagonal systems and (4,6)-fullerenes are equal to their Fries numbers. The previous paper also characterized the hexagonal systems H with af(H) = 2.

The inner dual graph T(H) of a hexagonal system H is the plane graph whose vertices

consist of the centers of all hexagons of H, and two centers are joined by a segment as an edge of T(H) if and only if the corresponding two hexagons have a common edge in H. Obviously, T(H) is a triangulation graph.



Figure 1. A general CHS $H(a_1, a_2, a_3, a_4, a_5, a_6)$.

A hexagonal systems is called *convex* if the boundary of its inner dual graph bounds a convex point set in the plane (see Fig. 1). Next a convex hexagonal systems will be abbreviated as CHS. Cyvin [5] showed that a CHS has a perfect matching if and only if it has equal opposite sides, that is, $a_1 = a_4$, $a_2 = a_5$ and $a_3 = a_6$. Such a CHS can be represented by $H(a_1, a_2, a_3)$. For example, see Fig. 2; some special cases are linear hexagonal chain $H(1, 1, a_3)$ and parallelogram HS $H(1, a_2, a_3)$ (see Fig. 2(a) and 2(b)).



Figure 2. Various examples for CHS $H(a_1, a_2, a_3)$.

In the following, we will discuss those CHS with at least one perfect matching. For $H(a_1, a_2, a_3)$, Cyvin [5] and Bodroža et al. [3] gave and proved a formula for the number of Kekulé structures. Zhang [22] got an expression of the Clar number, thus the maximum forcing number. For benzenoid parallelogram $H(1, a_2, a_3)$, Zhang and Li [21] and Li [12] showed that its minimum forcing numer and anti-forcing number are both equal to 1 respectively, and Zhao and Zhang [25, 26] derived explicit expressions of its forcing polynomial and anti-forcing polynomials by generating functions. For the definition of the forcing polynomial of a graph, see [24]. Zhu [29] showed that f(H(2, 2, n)) = 2 for all $n \geq 2$, and f(H(3, 3, n)) = 3 for all $n \geq 3$.

In this paper, we obtain the minimum forcing and anti-forcing numbers of $H(a_1, a_2, a_3)$ by applying the perfect path systems.

Theorem 1.1. Let $H(a_1, a_2, a_3)$ be a CHS. Then $f(H(a_1, a_2, a_3)) = af(H(a_1, a_2, a_3)) = \min\{a_1, a_2, a_3\}.$

2 Preliminary

In this section, we will introduce some notations and definitions. Let H be a hexagonal system embedded in the plane with some edges vertical. A *peak* (*valley*) of H is a vertex, all neighbors of which are below (above) it. Surely, all peaks and valleys have degree 2. Since H is a bipartite graph, its vertices can be colored white and black so that each pair of adjacent vertices receives distinct colors. We make a convention that all peaks (resp. valleys) are black (resp. white).

A monotone path system of H is a set of disjoint down paths of H, in which each path issues from a peak and ends at a valley. A *perfect path system* of H is a monotone path system covering all peaks and valleys of H.

Let M be a perfect matching of a graph G. We call a cycle (or path) of G is Malternating, if its edges appear alternately in M or not. Sachs [15] gave a one-to-one correspondence between the perfect matchings and the perfect path systems of a hexagonal system.

Theorem 2.1. [15] Let H be a hexagonal system. Then H has a perfect matching if and only if it has a perfect path system. Further their one-to-one correspondence is given as follows (see Fig. 3).

- (i) For a perfect path system *P*, all non-vertical edges in and vertical edges not in the paths of *P* form a perfect matching of *H*.
- (ii) If a perfect matching M is given, then we can delete all oblique edges not in M and all vertical edges in M together with their end vertices to obtain a perfect path system, where every path is an M-alternating path.



Figure 3. A perfect path system of H corresponding to a perfect matching and consisting of monotone paths from p_i to v_i , $1 \le i \le 4$.

Let M be a perfect matching of an HS H. Let D(H, M) be a digraph obtained from H by directing every edge of M from black to white end-vertices, and directing each edge in $E(H)\backslash M$ from white to black end-vertices. Then there is a natural one-to-one corresponding between the directed cycles (resp. paths) and M-alternating cycles (resp. paths). Obviously, all arcs (directed edges) in D(H, M) can be divided into six directions as shown in Fig. 4.



Figure 4. Six edge directions of a hexagonal system H.

An HS H has three plane drawings H_1 (the same as H), H_2 and H_3 such that they always have vertical edges and all peaks and valleys are black and white respectively. More precisely, rotating H 120 degrees counterclockwise and clockwise, we get H_2 and H_3 respectively. By Theorem 2.1, H_i has a perfect path system $\mathcal{P}_i(M)$ corresponding to M for each $i \in \{1, 2, 3\}$. When H_2 and H_3 are returned to H, perfect path systems $\mathcal{P}_i(M)$ for i = 2 and 3 are imposed into H. In this point of view, H has three perfect path systems $\mathcal{P}_i(M)$ for i = 1, 2 and 3. For each path Q_i in $\mathcal{P}_i(M)$, its orientation in D(H, M)is a directed path. For convenience, we also use Q_i to represent its orientation; for an edge e = xy in H, we also use e = xy to represent an arc (i.e. directed edge) in D(H, M)such that x and y are tail and head respectively.

In reference to the same H, we have the following three simple observations about each (directed or *M*-alternating) path in $\mathcal{P}_i(M)$ for i = 1, 2 and 3.

Observation 2.2. For each directed path $Q_1 \in \mathcal{P}_1(M)$, all arcs of Q_1 have directions \vec{i} , \vec{k} and $-\vec{j}$, and Q_1 is a down path.

Observation 2.3. For each directed path $Q_2 \in \mathcal{P}_2(M)$, all arcs of Q_2 have directions \vec{i} , \vec{j} and $-\vec{k}$. So all vertices of Q_2 are from right to left and the black (resp. white) vertices are from bottom to top along Q_2 .

Observation 2.4. For each directed path $Q_3 \in \mathcal{P}_3(M)$, all arcs of Q_3 have directions \vec{j} , \vec{k} and $-\vec{i}$. So all vertices of Q_3 are from left to right and the black (resp. white) vertices are from bottom to top along Q_3 .

To prove our Theorem 1.1 we present some basic results on general graphs and topology. The following well-known Jordan curve theorem in the plane will be used repeatedly.

Theorem 2.5. [13] (Jordan curve theorem) Any simple closed curve D divides the points of the plane not on D into two distinct domains (with no points in common) of which D is the common boundary.

Equivalent definitions for a forcing set and an anti-forcing set of a perfect matching of a graph are described as follows.

Lemma 2.6. [1] Let M be a perfect matching of a graph G. Then an edge subset of M is a forcing set of M if and only if each M-alternating cycle intersects S.

Lemma 2.7. [11] Let G be a graph with a perfect matching M. Then $S \subseteq E(G) \setminus M$ is an anti-forcing set of M if and only if every M-alternating cycle intersects S.

The following gives relations between forcing and anti-forcing numbers of a graph. **Lemma 2.8.** [11] Let G be a graph with a perfect matching M. Then $f(G, M) \leq af(G, M) \leq (\Delta - 1)f(G, M)$, and thus $f(G) \leq af(G)$.

3 Proof of Theorem 1.1

Our proof to Theorem 1.1 will be divided into three steps in this section. First we apply the three perfect path systems of $H(a_1, a_2, a_3)$ to get that the minimum forcing number of $H(a_1, a_2, a_3)$ has at least min $\{a_1, a_2, a_3\}$. Then, by constructing a 3-coordinate system of $H(a_1, a_2, a_3)$, we can find a special perfect matching M of it such that $af(H(a_1, a_2, a_3), M)$ $\leq \min\{a_1, a_2, a_3\}$. That implies that $H(a_1, a_2, a_3)$ has the minimum anti-forcing number at most min $\{a_1, a_2, a_3\}$. Finally, by combining a known result [11] stating that the minimum forcing number of a graph is equal to or less than its minimum anti-forcing number (see also Lemma 2.8) with the previous facts, we can directly get that the minimum forcing and anti-forcing numbers of $H(a_1, a_2, a_3)$ are both equal to min $\{a_1, a_2, a_3\}$.

Lemma 3.1. Let M be a perfect matching of a hexagonal system H. If three paths $Q_i \in \mathcal{P}_i(M), i \in \{1, 2, 3\}$, pairwise intersect, then there is an M-alternating cycle of H formed by three subpaths on the three paths Q_i .

Proof. Since Q_i and Q_j are both *M*-alternating paths, $1 \le i < j \le 3$, they have at least one common edge belonging to *M*. Let $e_1 = u_1v_1$ be a common edge of Q_1 and Q_2 so that $e_1 \in M$. By Observations 2.2 and 2.3, e_1 is of direction \vec{i} .

Claim 1. Any pair of Q_1 , Q_2 and Q_3 contains a unique common edge.

Proof. We only consider Q_1 and Q_2 . The other cases are similar. Let $e = wu_1$ and $e' = v_1w'$ be two edges of Q_1 adjacent to e_1 , if they exist. Then e and e' are both with direction $-\vec{j}$ and do not belong to Q_2 by Observation 2.2, which implies that w and w' lie on the right and on the left sides along the direction of Q_2 , respectively. We can deduce that Q_1 goes through Q_2 always from the right to the left of Q_2 as shown in Fig. 5(a). That implies that the subpath of Q_1 from the peak to u_1 lies entirely on the right side of Q_2 . Otherwise, we only need to consider the case that the latter does not holds. That is, the directed subpath Q'_1 of Q_1 from w' to the end enters in the right side of Q_2 . By the Jordan curve theorem, Q'_1 intersects Q_2 at least one point, and Q'_1 goes through Q_2 from the left to the right of Q_2 , a contradiction. So Q_1 and Q_2 have exactly one common edge.



(c) Case 2

(d) An *M*-alternating cycle bounded by Q_1, Q_2 and Q_3 .

Figure 5. Illustration for the proof of Lemma 3.1.

By Claim 1, e_1 is the common edge of Q_1 and Q_2 , and let $e_2 = u_2v_2$ (resp. $e_3 = u_3v_3$) be the common edge of Q_2 and Q_3 (resp. Q_1 and Q_3). By Observations 2.2 to 2.4, e_2 is of direction \vec{j} , and e_3 is of direction \vec{k} . So e_1, e_2 and e_3 are pairwise different edges in M. So we know that the u_i and v_i receive black and white, $1 \le i \le 3$, respectively.

Since e_i and e_{i-1} are both in Q_i , we can get a directed sub-path Q''_i of Q_i , connecting but not containing e_i and e_{i-1} , where i = 1, 2, 3, and $e_0 = e_3$. Then each Q''_i is an *M*alternating path whose both end edges are not in *M*. So the first and final vertices of Q''_i respectively belong to $\{v_1, v_2, v_3\}$ and $\{u_1, u_2, u_3\}$. By Claim 1, any pair of Q''_1 , Q''_2 and Q''_3 has no common edges. Further we have the following claim.

Claim 2. Any pair of Q_1'' , Q_2'' and Q_3'' has no common vertices.

Proof. For two vertices x and y in Q_i such that x appears before y along Q_i , let $Q_i[x, y]$ denote the (directed) subpath of Q_i from x to y. There are two cases to distinguish.

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Case 1. e_1 comes before e_3 along the path Q_1 (see Fig. 5(b)).

Then v_1 lies above u_3 and $Q''_1 = Q_1[v_1, u_3]$. It suffices to prove that e_i is before e_{i-1} along the path Q_i for i = 2, 3, that is, $Q''_2 = Q_1[v_2, u_1]$ and $Q''_3 = Q_3[v_3, u_2]$. Since $e_1 = u_1v_1$ is of direction \vec{i} , u_1 is above v_1 . The entire subpath of Q_2 from v_1 lies on the right of Q_1 by Claim 1 and above v_1 by Observation 2.3. Since $e_3 = u_3v_3$ is of direction \vec{k} , u_3 is above v_3 . By Claim 1 and Observation 2.4, the entire subpath of Q_3 to u_3 lies on the right of Q_1 and below u_3 . So the subpath of Q_2 from v_1 and the subpath of Q_3 to u_3 are disjoint. Further, since Q_2 and Q_3 intersect at edge e_2 , e_2 is the common edge of the subpath of Q_2 from the initial vertex to u_1 and the subpath of Q_3 from v_3 to the final vertex. That implies that e_3 appears before e_2 along the path Q_3 and Q_2 first passes through e_2 before e_1 .

Case 2. e_3 comes before e_1 along the path Q_1 (see Fig. 5(c)).

Then v_3 is above u_1 and $Q''_1 = Q_1[v_3, u_1]$. Similar to Case 1, we just need to show that $Q''_2 = Q_2[v_1, u_2]$ and $Q''_3 = Q_3[v_2, u_3]$. Since u_1 is above v_1 . By Claim 1 and Observation 2.3, the entire subpath of Q_2 to u_1 lies on the left of Q_1 and below u_1 . Since u_3 is above v_3 . The entire subpath of Q_3 from v_3 lies on the left of Q_1 and above v_3 by Claim 1 and Observation 2.4. Thus, the subpath of Q_2 to u_1 and the subpath of Q_3 from v_3 are disjoint. That is, e_2 is the common edge of the subpath of Q_2 from v_1 to the final vertex and the subpath of Q_3 from the initial vertex to u_3 . That implies that $Q''_2 = Q_2[v_1, u_2]$ and $Q''_3 = Q_3[v_2, u_3]$.

From Claim 2, we know that the end vertices of Q_1'', Q_2'' and Q_3'' are different. So we can find a directed cycle $e_1 \cup Q_1'' \cup e_3 \cup Q_3'' \cup e_2 \cup Q_2''$, also an *M*-alternating cycle.

Let $H(a_1, a_2, a_3)$ be a CHS with a perfect matching M. Without loss of generality, we can assume that the number of peaks and valleys of H_i are both a_i for i = 1, 2, 3. By Theorem 2.1 the number of paths in a perfect path system of H_i is a_i , that is $|\mathcal{P}_i(M)| = a_i$.

Lemma 3.2. Let $H(a_1, a_2, a_3)$ be a CHS with a perfect matching and $a_1 \leq a_2 \leq a_3$. Then $f(H(a_1, a_2, a_3)) \geq a_1$.

Proof. Assume to the contrary that $f(H(a_1, a_2, a_3)) < a_1$. Let M be a perfect matching of $H(a_1, a_2, a_3)$ with the minimum forcing number. Then M has a forcing set S with cardinality smaller than a_1 . By Theorem 2.1 (ii), for each $i \in \{1, 2, 3\}, \mathcal{P}_i(M)$ consists of a_i disjoint directed paths of D(H, M). So there are three directed paths $Q_1 \in \mathcal{P}_1(M)$, $Q_2 \in \mathcal{P}_2(M)$ and $Q_3 \in \mathcal{P}_3(M)$ which each contains no edges of S.

For each pair $\{i, j\} \subseteq \{1, 2, 3\}$, since both Q_i and Q_j are inside the boundary of H(a closed polygon), and their end vertices cross along the boundary, by Jordan curve theorem, Q_i and Q_j intersect at least one point. By Lemma 3.1, Q_1 , Q_2 and Q_3 form an M-alternating cycle containing no edges of S, contradicting that S is a forcing set of Mby Lemma 2.6. Thus, we have $f(H(a_1, a_2, a_3)) \ge a_1$.

For any edge e of a graph G, we call it a 2-2 edge, if both end vertices of e have degree 2 in G. Obviously, $H(a_1, a_2, a_3)$ has exactly six 2-2 edges, and they are all on the border of $H(a_1, a_2, a_3)$. We label the six 2-2 edges by $e_1, e_2, e_3, e_4, e_5, e_6$ in clockwise order along the border of $H(a_1, a_2, a_3)$. Moreover, e_i and e_{i+3} are parallel, where i = 1, 2, 3.



Figure 6. Illustration for the proof of Lemma 3.4: three anti-forcing sets of a perfect matching of $H(a_1, a_2, a_3)$, where $a_1 = 2$, $a_2 = 4$ and $a_3 = 6$.

Lemma 3.3. [20, 23] Let H be a hexagonal system with a perfect matching M. Then for any M-alternating cycle C of H, there is an M-alternating hexagon in C with its interior. **Lemma 3.4.** Let $H(a_1, a_2, a_3)$ be a CHS with a perfect matching and $a_1 \leq a_2 \leq a_3$. Then $af(H(a_1, a_2, a_3)) \leq a_1$.

Proof. Take rays OA, OB and OC that are perpendicular bisectors of 2-2 edges e_1 , e_3 and e_5 respectively so that they intersect at the center O of some hexagon h (see Fig. 6). Then O - ABC can be seen as a 3-coordinate system of H, which divides the plane into three areas AOB, BOC and COA. Let L_A (resp. L_B and L_C) be the set of edges

of *H* intersecting *OA* (resp. *OB* and *OC*). Then the cardinalities of L_A , L_B and L_C respectively are a_1 , a_2 and a_3 .

Let M be a perfect matching of H such that M does not contain an edge of L_A , L_B and L_C , and all edges of M in anyone of the three areas are parallel to each other, see Fig. 6. Then H contains exactly one M-alternating hexagon h whose center is O. For any M-alternating cycle C of H, by Lemma 3.3, h lies in the C with its interior, so O is a point in the interior of C. So each one of L_A , L_B and L_C intersects C. In addition, L_A , L_B and L_C are subsets of $E(H)\backslash M$. By Lemma 2.7, L_A , L_B and L_C are all anti-forcing sets of M. Hence we have that $af(H) \leq af(H, M) \leq a_1$.

Proof of Theorem 1.1. Without loss of generality, suppose that $a_1 = \min\{a_1, a_2, a_3\}$. Combining Lemmas 3.2, 3.4 and 2.8, we have the following inequalities.

$$a_1 \le f(H(a_1, a_2, a_3)) \le af(H(a_1, a_2, a_3)) \le a_1,$$

which implies that all the above equalities hold.

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