

A New Protein Domain Assignment Algorithm Based on the Dominating Set of a Graph

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Abstract

Assignment of structural domains in complex protein structures is an important task in bioinformatics researches. As the number of known protein structures grows rapidly, the need for automatic methods for determining protein domains based on the proteins tree-dimensional structure becomes more desirable. In this paper, we introduce a new domain decomposition algorithm which is based on the dominating set of the graph representation of a protein. To evaluate our method, we compare our results with the other computational methods on a commonly used benchmark of 55 proteins. It is shown that the performance of our algorithm is better than the other automatic methods.

Introduction

Proteins can be considered as a set of several structural domains. Each domain has a stable structure and can fold independently of the rest of the protein [1–3]. Structural domains are compact and should have a hydrophobic core. Each of these semi-independent units has a specific function [4].

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Structural domains are the basic components of the proteins. They should not necessarily be continuous in the amino acid sequence and may consist of non-sequential segments [5, 6]. The assignment of structural domains is an important task in the classification of the proteins based on their three-dimensional structure [7, 8], understanding the proteins folding, function and evolution [9]. The concept of assigning protein domains has been proposed by Wetlauffer [6], Rossmann and Liljas [10] in 1970. Domain decomposition can be done manually by human experts. There are several classifications of the protein structures based on structural domains like SCOP [7] and CATH [8]. SCOP classifications rely mainly on human experts. CATH uses both automatic methods and human experts' opinion for the classification of the protein structures. Due to the exponential rate of growth in the identification of the protein structures, the need for automatic methods for determining protein domains are required [11]. There are several automatic algorithms such as NCBI [12], DomainParser [13], PDP [14], PUU [15], DDomain [16], DHcl [17] and Dodis [18]. The computational approaches of these methods are different but they mainly focus on the fact that the residue contacts of amino acids within a domain are denser than between domains [19]. In this paper, we introduce a novel algorithm for determining protein domains, using the dominating set of the graph representation of a protein.

Method

A graph is usually shown by $G = (V, E)$ where V is a finite set of nodes and E is a finite set of edges, which are 2-element subsets of V . For constructing the graph of a protein, each amino acid residue of the protein is considered as a node of a graph. The edges of this graph are generated from the structural coordinates of the amino acid residues [20] that are obtained from the PDB (Protein Data Bank) [21]. Two nodes are connected by an edge if the distance between the C^α atoms of their corresponding amino acid residues is 4\AA or less, following the definition of Holm and Sander [9].

A dominating set for a graph $G = (V, E)$ is a subset D of V such that every vertex not in D is a neighbor of at least one member of D . For example given the graph G shown in Figure 1, $D = \{3, 5\}$ is a dominating set for G .

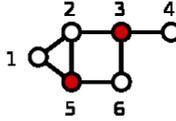


Figure 1. The set $\{3, 5\}$ denotes a dominating set for graph G .

For assigning the protein domains, we first construct a dominating set for the protein’s graph. Let P be a protein with m amino acids and $G_p = (V_p, E_p)$ shows its representative graph. Finding the minimum dominating set in a graph is an NP-hard problem [22], so we use a greedy approach to obtain a dominating set D for the graph G_p . A dominating set for the protein 1A8Y is shown in Figure 2. The graph of this protein consists of 347 nodes while its dominating set has 60 nodes.

• 4 → 3 8 6	• 193 → 150 151 152 153 159 160 192 194 195 211 212 213
• 9 → 7 8 10 11 12 60 63	• 200 → 143 144 145 146 198 199 201 202 228 229 230 231
• 22 → 17 18 19 20 21 23 24 25 26	• 205 → 196 197 198 199 200 203 204 206 207
• 23 → 18 19 20 21 22 24 25 26 27 63	• 213 → 159 162 163 193 194 195 208 309 310 311 212 214 215 217 218
• 25 → 22 23 24 26 27 28 29	• 214 → 159 162 163 212 213 215 216 217 218
• 37 → 35 36 38 74 75 76 77 84 89 91 92 93 94 95 96	• 223 → 219 220 221 222 224 225 226
• 48 → 41 42 43 44 45 46 47 49 50 51 52	• 229 → 200 227 228 230 231 273 276 377 380 282 283 284
• 50 → 46 47 48 49 51 52 53 54	• 231 → 144 145 200 201 227 228 229 230 232 233 241 283 284 285 286
• 59 → 13 14 15 50 56 57 58 60 61 62 63	• 235 → 108 109 114 233 234 236 237 238 291 292
• 60 → 9 56 57 58 59 61 62 63 64 112	• 236 → 106 107 108 109 114 115 116 118 234 235 237
• 62 → 14 15 58 59 60 61 62 64 65 66 70 71 72	• 241 → 231 232 233 237 238 239 342 343 344 345 285
• 63 → 9 14 20 60 61 62 64 65 66	• 245 → 241 242 243 244 246 247 248 249 250 283 317
• 65 → 61 62 63 64 66 67 68 70 116	• 255 → 253 254 256 257 262 263 265 266 286 287 288 289 309 310
• 71 → 15 16 30 31 34 62 66 68 70 72 73	• 263 → 255 257 258 260 261 302 304 255 256 267
• 76 → 37 38 39 40 74 75 77 78 79 80	• 272 → 268 269 270 271 273 274 275 276 336 337 339 340
• 84 → 37 75 77 80 81 82 83 85 86 87 88 89 90	• 280 → 229 277 278 279 281 282
• 89 → 31 32 97 98 100 101 102 103 104 122 123	• 286 → 230 231 232 235 252 253 254 255 284 285 287 288
• 103 → 98 97 98 99 101 102 104 105	• 288 → 233 254 255 256 257 286 307 389 390 291 292
• 114 → 108 109 110 111 112 113 115 116 117 118 235 236	• 289 → 254 255 256 287 288 290 291 292 293 296 307
• 122 → 99 118 119 120 121 123 124 125 171 172	• 292 → 235 288 289 290 291 293 294 295 296
• 136 → 130 132 133 134 135 137 138 139 140 184	• 296 → 289 292 293 294 295 297 289 300 307
• 138 → 134 135 136 137 139 140 141 142	• 300 → 296 297 289 299 301 302 303 304 305 306 307
• 152 → 150 151 153 154 155 157 159 160 178 179 180 192 193	• 306 → 300 304 305 307 308 309
• 157 → 152 153 154 155 156 158 159 160 161	• 315 → 249 250 251 252 313 314 316 317 320 321 322
• 159 → 152 157 158 160 161 162 163 163 212 213 214	• 316 → 248 249 314 315 317 318 319 320 321 322
• 167 → 126 127 163 164 165 166 168 169 170 171 174 175 176	• 325 → 311 312 313 323 324 326 342
• 175 → 126 127 128 147 148 167 171 172 173 174 176 177	• 339 → 272 334 335 336 337 338 340 341 342 343
• 179 → 129 130 131 149 150 151 152 178 180 181 184 185	• 340 → 272 336 337 338 339 341 342 343 344
• 184 → 130 131 132 133 136 179 181 182 183 185 186 187 188	• 342 → 325 338 339 340 341 343 344 345 346
• 185 → 151 179 180 181 182 183 184 186 187 188 189 190 191	• 343 → 339 340 341 342 344 345 346 347

Figure 2. A dominating set of the protein 1A8Y.

Next we construct a matrix for the obtained dominating set, $D = \{x_1, x_2, \dots, x_n\}$. We define a matrix $DS = [DS_{i,j}]$ by:

$$DS_{i,j} = \frac{|N(x_i) \cap N(x_j)|}{|N(x_i) \cup N(x_j)|}$$

where $N(x_i)$ denotes the set of the neighbors of the node x_i in G_p .

Figure 3 shows the matrix DS for the dominating set of the protein 1A8Y. This protein has three domains which are shown by different colors in Figure 4. Its initial domains are also shown by different colors in the DS matrix. The entries of the colored parts of the DS matrix are almost none zero, while the rate of none zero elements in the white parts is small. The decomposition of the domains of this protein is (3-126), (127-228) and (229-347).

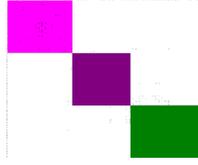


Figure 3. The matrix DS of the dominating set of the protein 1A8Y.

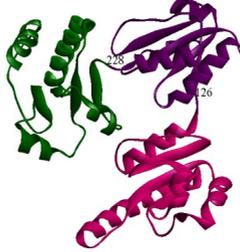


Figure 4. A solid ribbon diagram showing the three domains of the protein 1A8Y.

For determining and merging the initial clusters, first we define the distance matrix $DIS = [d_{i,j}]$ from the matrix DS as follows:

$$d_{i,j} = \frac{\sum_k |DS_{i,k} - DS_{j,k}|}{n}$$

The members of the dominating set are considered as initial clusters and are merged based on this distance matrix and the neighbor-joining algorithm [23]. For this purpose, first the array U of size n is obtained from the matrix DIS by:

$$U_i = \frac{1}{n-2} \sum_{i \neq k} d_{i,k}$$

Then the matrix M is constructed from U and DIS :

$$M_{i,j} = d_{i,j} - U_i - U_j$$

We define θ as:

$$\theta = \frac{\min[m_{i,j}] + \max[m_{i,j}]}{3}$$

For merging the clusters, the minimum entry of M , $M_{x,y}$, is selected and the clusters x and y are merged together. Then the distance matrix DIS is updated by changing the row corresponding to the cluster x as:

$$d_{x,k} = \frac{d_{x,k} + d_{y,k} - d_{x,y}}{2}$$

and removing the row corresponding to y . The matrixes U and M are then computed from the matrix DIS in each step. This procedure is repeated until $M_{x,y}$ is less than θ .

In the next step, obtained clusters are merged based on their inter and intra densities; with respect to the fact that the residue interactions are denser within domains than between domains [19]. The density of the cluster C_i is computed by:

$$density(C_i) = \frac{|E(C_i)|}{|C_i|}$$

where $|E(C_i)|$ denotes the number of edges between the nodes of C_i . The intra-residue interactions of a cluster, which is the result of merging two clusters C_i and C_j is defined as:

$$intradensity(C_i, C_j) = \frac{|E(C_i \cup C_j)|}{|C_i \cup C_j|}.$$

The inter density between two clusters C_i and C_j is computed by:

$$interdensity(C_i, C_j) = \frac{|E(I(C_i \cup C_j))|}{|C_i \cup C_j|}$$

where $|E(I(C_i \cup C_j))|$ denotes the set of edges with one end in C_i and the other end in C_j . We define the total density of the two clusters C_i and C_j as:

$$totaldensity(C_i, C_j) = intradensity(C_i, C_j) - \frac{(density(C_i) + density(C_j))}{2} + interdensity(C_i, C_j).$$

Two clusters C_i and C_j with the maximum total density are repeatedly merged together until the number of clusters become less than η .

Next unassigned vertices are determined and merged with the existing clusters based on their neighbors in each cluster.

In the next step, we assign a pattern to each cluster. Let $V_P = \{v_1, v_2, \dots, v_m\}$, we define the $m \times m$ matrix NA as:

$$NA_{i,j} = |N(v_i) \cap N(v_j)|.$$

For a cluster C , the pattern $P(C)$ is defined by:

$$P(C) = \sum_{x \in C} r_x$$

where r_x is the row corresponds to the node x in NA . Then the similarity score $S(C, D)$ between two clusters, C and D , is defined by:

$$S(C, D) = \frac{|\{k|k \in C \text{ and } P(D)_k \neq 0\}|}{|C|}.$$

Two clusters X and Y with the maximum similarity score are repeatedly merged until $S(X, Y)$ become less than a threshold δ .

Threshold determination

The thresholds that have been used in this algorithm are determined using a training set consisting of 50 proteins selected from a set of 135 proteins in the Balanced Domain Benchmark-3 of the pDomain resource introduced in [4]. This database is available at <http://www.pdomains.sdsc.edu>. Both expert methods, CATH and SCOP, agree on the domain decomposition of these 50 proteins, which are selected as the training set.

The obtained values for the parameters are: $\eta = 10$ and $\delta = 45$. The minimum size of a domain is considered to be 32 residues in our algorithm.

Results and Discussion

The algorithm is applied to a frequently used benchmark consisting of 55 proteins introduced by Jones et al. [24]. A domain assignment is considered correct if the number of domains is the same as the assignment by the experts and the amino acid assignment of the domains is at least 85% in agreement with the experts' opinion [24]. In this paper, the domain decomposition of the automatic methods is compared with the assignments by the human experts, CATH or SCOP, similar to [4]. Using the above definition, the domain decomposition of each method is considered correct if it is consistent with the domain assignment of CATH or SCOP. It is noticeable that even the manual assignments of the protein domains, are sometimes different for the same proteins; since there is not a precise definition of protein domains [25–27]. This could also be the result of considering the function and evolutionary information of proteins in the domain decompositions by experts [28].

Our method correctly assigns 96.3% of the 55 proteins (Table 1). To compare our results with the assignments of other automatic domain assignment methods, we use dConsensus. dConsensus is a web resource which is available at <http://pdomains.sdsc.edu/dConsensus> [4] and displays the results of domain decompositions from multiple algorithmic methods. Using this software the results of six automatic domain assignment algorithms is calculated.

According to these results, the correct assignments by PDP, DomainParser, NCBI and PUU are 92.7%, 85.5%, 89% and 76.4% respectively. DHcL and DDomain run only on 41 and 50 proteins and their results are 70.7% and 84%.

Table 1. Protein PDB codes of 55 proteins, residue ranges of domains assigned by CATH, SCOP and our algorithm (fragments of domains are separated by ‘;’ and ‘/’ is used to separate domains).

Protein PDB ID	CATH	SCOP	Our Algorithm
8acna	2-202/ 203-315/ 316-490/ 534-754	2-528/ 529-754	2-528/ 529-754
3pmga	1-197/ 198-300/ 301-400/ 401-561	1-190/ 191-303/ 304-420/ 421-561	1-188/ 189-300/ 301-420/ 421-561
1phha	1-72, 96-180, 269-351/ 73- 95, 181-268, 352-388	1-173, 276-394/ 174-275	1-180, 267-394/ 181-266
3grsa	18-160, 290-365/ 161-289/ 366-478	18-165, 291-363/ 166-290/ 364-478	18-150, 290-363/ 151-289/ 364-478
1atna	5-35, 72-135, 338-373/ 36- 69/ 137-182, 272-333/ 183- 268	2-147/ 148-373	1-179, 247-372/ 180-273
1ezma	1-152/ 153-298	1-301	1-132/ 133-298
1fnba	19-151/ 152-314	19-154/ 155-314	19-163/ 164-314
1gpba	19-485 , 813-836/ 486-812	1-842	19-484, 813-841/ 485-812
1lapa	1-165/ 166-483	1-159/ 160-484	1-170/ 171-484
1pfka	1-142, 257-303/ 143-252, 304-319	1-320	1-137, 257-302/ 138-256, 303-319
1ppna	1-212	1-212	1-212
1rhda	1-156/ 157-293	1-149/ 150-293	1-151/ 152-293
1sgta	1-12, 97-210/ 13-96, 211- 223	1-223	1-223
1vsga	1-33, 86-255/ 34-85 256- 362	1-364	1-24, 86-253/ 42-85, 254- 362
1bksa	1-267	1-268	1-267
2cypa	4-144, 266-294/ 145-265	1-294	2-140, 255-294/ 141-254
2hada	1-310	1-310	1-310
3cd4a	1-98/ 99-173	1-97/ 98-178	1-98/ 99-178
1g6na	10-138/ 139-207	8-138/ 139-207	7-137/ 138-206
3pgka	2-187/ 194-402	1-416	1-200/ 201-415
4gcra	1-83/ 84-174	1-85/ 86-174	1-80/ 81-174

5fbpa	7-199/ 200-334	1-335	6-200/ 201-334
8adha	1-178, 318-374/ 179-317	1-163, 340-374/ 164-339	1-189, 322-374/ 190-321
8atca	1-133, 292-310/ 134-291	1-150/ 151-310	1-134, 285-310/ 135-284
8atcb	8-100/ 101-153	8-100/ 101-153	8-97/ 101-153
2acea	4-535	1-537	4-317/ 318-535
2buka	13-196	13-196	26-195
2aaka	1-150	1-152	1-150
1bbha	1-131	1-131	1-131
1bbpa	1-173	1-173	1-173
1brda	8-226	1-248	8-226
1fxia	1-96	1-96	1-96
1gkya	2-33, 94-187/ 34-93	1-187	2-33, 82-186/ 34-81
2gmfa	4-124	1-127	4-124
1gmpa	1-96	1-96	1-96
1goxa	2-360	1-370	2-360
1ofva	1-169	1-169	1-169
1pypa	1-281	1-285	1-281
1rbpa	1-175	1-182	1-175
1rcba	1-129	1-129	1-129
1rvea	2-245	1-245	2-245
1snca	7-141	1-149	7-141
1tiea	1-170	1-172	1-170
1tlka	33-135	1-154	33-135
1ulaa	1-289	1-289	1-289
1bkbsb	9-53, 87-205/ 54-86, 206-391	1-397	3-394
2azaa	1-129	1-129	1-129
2ceya	1-306	1-306	1-306
2m2a	1-155	1-155	1-155
2tmvp	1-154	1-158	1-154
3chya	1-128	1-128	1-128
3claa	1-213	1-213	1-213
3dfra	1-213	1-213	1-213
4blma	1-162	1-162	1-162
5p21a	1-166	1-166	1-166

The proteins that are decomposed incorrectly by our method are 1atna and 2acea (Figure 5).

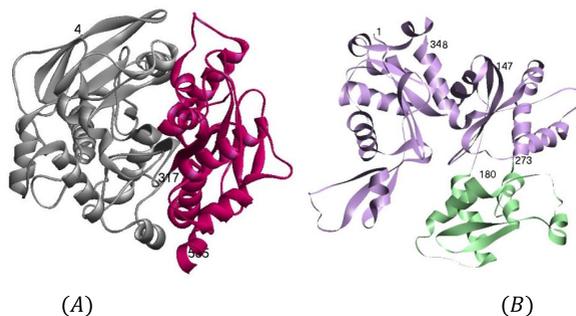


Figure 5. Domain decompositions of the proteins 2acea and 1atna which is obtained by our algorithm. Different domains are shown by different colors. (A) 2acea (4-317/ 318-535). (B) 1atna (1-179, 247-372/ 180-273).

The protein 2acea is considered as a one domain protein by the experts, but our algorithm assigns two domains to this protein (Figure 5(A)). Among automatic methods, DomainParser and DDomain consider this protein as a single domain protein. DHcL assigns two domains for this protein which is similar to our algorithm (Table 2).

Table 2. Residue ranges of domains assigned by different methods for protein PDB code 2acea (fragments of domains are separated by ‘,’ and ‘/’ is used to separate domains).

2acea					
SCOP	1-537	CATH	4-535	OUR Algorithm	4-317/ 318-535
pdp	4-315/ 332-394, 526-535/ 316-331 , 395-525	DomainParser	1-537	NCBI	1-230, 301-326, 415-516/ 231-300/ 327-414, 517-537
puu	1-233, 281-332, 396-508/ 234-280/ 333-395	DDomain	4-535	DHcL	4-315/ 316-535

For the protein 1atna, expert methods give different domain decompositions. SCOP considers this protein as a two-domain protein while CATH assigns four domains for this

protein. Our algorithm considers two domains for this protein (Figure 5(B)) similar to SCOP but the fragments of our domains are inconsistent with the assignment by SCOP. Only pdp considers four domains for this protein similar to the CATH assignment (Table 3). Domain decomposition by DomainParser is also similar to our assignment.

Table 3. Residue ranges of domains assigned by different methods for protein PDB code 1atna (fragments of domains are separated by ',' and '/' is used to separate domains).

1atna					
SCOP	2-147/ 148-373	CATH	5-35, 72-135, 338-373/ 36-69/ 137-182, 272- 333/ 183-268	OUR Algorithm	1-179, 247- 372/ 180-273
pdp	2-34, 70-138, 340-373/ 35- 69/ 139-185, 261-339/ 186- 260	DomainParser	2-148, 338-373/ 149-337	NCBI	1-137, 353- 372/ 138-182, 263-352/ 220- 262
puu	1-33, 69-141, 336-372/ 142- 179, 273-335/ 180-272	DDomain	2-103/ 104-373	DHcl	2-373

The above results show that our algorithm which is introduced in this paper performs better results compared to other automatic algorithms.

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