

A THEOREM ON GRAPH VALENCE SHELLS

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Abstract. Recently Randić [1] has proposed a set of new graph invariants: the graph valence shells. A valence shell $s(u)_k$ is the sum of valences (degrees) of all vertices placed at distance $k - 1$ from vertex u . The graph valence shell s_k is the sum of valence shells $s(u)_k$. Consider an acyclic graph T . It has been proved that $s_k = p_{k-1} + p_k$, where p_k denotes the number of paths of length k in T .

INTRODUCTION

Recently Randić [1] has proposed a set of new graph invariants: the graph valence shells. A valence shell $s(u)_k$ is the sum of valences (degrees) of the vertices placed at a constant distance $k - 1$ from a given vertex u . Randić has not given an explicit formula for $s(u)_k$, but the following expression is in full agreement with Randić's definition:

$$s(u)_k = \sum_i \delta(k - 1 - d_{u,i}) v_i \quad (1)$$

where $d_{u,i}$ denotes the distance between vertices u and i , and v_i denotes the valence of a vertex i , δ is a function which is equal to 1 if the argument is zero, and is zero in all other cases. δ picks out a vertex placed at distance $d_{u,i}$ from u . The summation has to be performed for all N vertices in T . Graph valence shell s_k can be obtained by summing expression (1) for all vertices u , and dividing the result by two:

$$s_k = \sum_u s(u)_k / 2 \quad (2)$$

The calculation of s_k can be illustrated by using a simple example (FIGURE 1). As an illustration obtain s_5 : The valence shell of vertex 1 is 2 (valence of vertex 5). Contribution of vertex 2 is 1 (valence of vertex 6), contribution of vertex 5 is $1 + 1 = 2$ (sum of valences of vertices 1 and 7), contribution of vertex 6 is 6 (sum of valences of vertices 2, 8 and 9). The contribution of vertex 7 is 2 (valence of vertex 5) and the contributions of vertices 8 and 9 are 1, each. Therefore $s_5 = (2 + 1 + 2 + 5 + 2 + 2)/2 = 7$. Graph invariants based on 'concentric shells' (or layers) have already been discussed [2-6] and graph valence shells happen to be a subclass of this more general approach.

Graph valence shells s_2 , s_3 and s_4 have been shown by Randić to form an appropriate set of indexes that can be used in structure-property relationships. More specifically Randić found that combinations of p_2 , p_3 , p_4 and s_2 , s_3 , s_4 , respectively, resulted in multiple regression equations with identical statistical parameters. (The boiling points of octanes were considered as the dependent variable.) This result could only have been obtained if p_2 , p_3 , p_4 and s_2 , s_3 , s_4 are linear transforms of each other. Actually Randić has observed that

$$p_1 + p_2 = s_2 \quad (3)$$

$$p_2 + p_3 = s_3 \quad (4)$$

$$p_3 + p_4 = s_4 \quad (5)$$

The aim of this paper is to prove that these relationships are valid, and that they are instances of a more general theorem:

Theorem: Let T be a tree and let p_k denote the number of distances of length k , while s_k denotes the k -th graph valence shell in T . Then

$$p_{k-1} + p_k = s_k \quad (6)$$

PROOF OF THE THEOREM

Let's first show that

$$p_1 + p_2 = s_2 \quad (7)$$

because the argumentation will reveal the relationship between s_2 and the Zagreb indices [7]. Vertices contributing to s_2 are the first neighbours of vertex u (eq. 2). Each vertex u has v_u first neighbours. For each first neighbour the contribution to s_2 is equal to v_u , therefore the total contribution of u is v_u^2 , and we obtain:

$$s_2 = \sum_u v_u^2/2 \quad (8)$$









Division by 2 has been introduced in accordance with Randić's definition [1]. Eq. (8) indicates the close relationship between s_2 and one of the Zagreb indices [7]. Eq. (8) may be rewritten:

$$s_2 = \sum_u v_u^2/2 = \sum_u v_u(v_u - 1)/2 + \sum_u v_u/2 \quad (9)$$

The second term on the right hand side of eq. (9) is equal to p_1 (the number of edges in T) because of the 'handshake' lemma. The first term is equal to the number of times a path of length two (l_2) can be placed on vertex u . Therefore the first term of eq. (9) is equal to p_2 and $p_1 + p_2 = s_2$.

In order to derive $p_2 + p_3 = s_3$, observe that all second neighbours of u have to be considered now. The second and third columns of Table 1 indicate how many times paths of length two (l_2) and of length three (l_3) can be placed upon the substructures k_{nm} [8,9]. Note that k_{nm} will denote the substructure and the number this structure appears in T . The fourth column of Table 1 can be obtained by inspection of the respective substructures k_{nm} . As an exercise consider substructure k_{13} : the endpoint will contribute (two times) to $s(u)_3$, and its contribution is 1 for each neighbour (free valence), therefore the total contribution of k_{13} is equal to 2. As a second example consider k_{22} : Each vertex will contribute 2 to $s(u)_3$, the sum of these contributions is equal to 4. In general the contribution of k_{nm} to $s(u)_3$ is equal to $(m-1)n + (n-1)m$.

TABLE 1. Substructures k_{mn} and their contributions to $s(u)_3$.
Endpoints are denoted by black dots.

Substructure	Number of times l_2 can be placed upon k_{mn}	Number of times l_3 can be placed upon k_{mn}	Contribution to $s(u)_3$
 k_{11}	0	0	0
 k_{12}	1	0	1
 k_{13}	2	0	2
 k_{14}	3	0	3
 k_{22}	2	1	4
 k_{23}	3	2	7
 k_{24}	4	3	10
 k_{33}	4	4	12
k_{mn}	$(m-1)+(n-1)$	$(m-1)x(n-1)$	$(m-1)n+(n-1)m$

The sum of substructures is equal to the number of edges [8,9]:

$$\sum_{mn} k_{mn} = p_1 \quad (10)$$

The number of paths of length two (p_2) is (Table 1):

$$\sum_{mn} k_{mn}(m+n-2) = 2p_2 \quad (11)$$

The factor of two on the right hand side appears because each vertex was considered two times in eq. (11). The number of paths of length three (p_3) is (Table 1):

$$\sum_{mn} k_{mn}(m-1)(n-1) = p_3 \quad (12)$$

Finally for s_3 we obtain (Table 1):

$$\sum_{mn} k_{mn}(2mn - m - n) = 2s_3 \quad (13)$$

Combining eqs. (11), and (12) we obtain:

$$2p_2 + 2p_3 = \sum_{mn} k_{mn}[(m + n - 2) + 2(m - 1)(n - 1)] = \sum_{mn} k_{mn}(2mn - m - n) = 2s_3 \quad (14)$$

and from this $p_2 + p_3 = s_3$ follows immediately.

Next prove that $p_{k-1} + p_k = s_k$. In order to obtain e.g. $p_3 + p_4 = s_4$, leave the free valences in TABLE 1 unaltered, and replace the edge with symbol $-$, where the square denotes a vertex of any degree. Replace symbols l_2 and l_3 by l_3 and l_4 , respectively, and modify the definition of substructure k_{mn} : k_{mn} will now denote the substructure (and the number of such substructures), which was obtained replacing the edge in k_{mn} by $-$. Then an analogous argumentation given for the $p_2 + p_3 = s_3$ case can be applied. Similarly if $-$ in turn is replaced by $-$ (etc.), the same reasoning can be applied. Therefore we obtain that $p_{k-1} + p_k = s_k$.

The theorem remains valid for p_0 and p_1 . Since p_0 (the number of paths of length zero) is equal to zero, and $p_1 = N - 1$ (the number of edges in T), $p_0 + p_1 = s_1$, because of the 'handshake' lemma.

DISCUSSION

Path-counts can easily be determined by inspection of the matrix D composed of the distances $d_{i,j}$. The number of times any number k appears in the upper right hand side of D is equal to p_k . Therefore all values of s_k can be obtained in a more simple manner, than by using eq. (1).

Eq. 6 can be expressed in matrix form:

$$Bp = s \quad (15)$$

where \mathbf{p} is vector containing path counts p_k ($k = 1, \dots, m$), and m denotes the value of the maximal distance in T . Similarly \mathbf{s} is a vector consisting values of the individual graph valence shells s_k ($k = 1, \dots, m$). The structure of matrix \mathbf{B} is the following in case $N = 5$:

$$\mathbf{B} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 \end{bmatrix}$$

Through inversion of \mathbf{B} we obtain:

$$\mathbf{B}^{-1} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ -1 & 1 & 0 & 0 & 0 \\ 1 & -1 & 1 & 0 & 0 \\ -1 & 1 & -1 & 1 & 0 \\ 1 & -1 & 1 & -1 & 1 \end{bmatrix}$$

And therefore:

$$\mathbf{p} = \mathbf{B}^{-1}\mathbf{s} \quad (16)$$

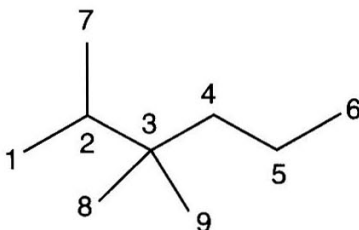


FIGURE 1. Hydrogen suppressed graph of 2,3,3-trimethyl-hexane.

The regular structure of B^{-1} will not change with other values of N , and we can write:

$$(B^{-1})_{ij} = \begin{cases} 0 & \text{if } i < j \\ (-1)^{i+j} & \text{if } i \geq j \end{cases}$$

and therefore expressing path counts in term of graph valence shells:

$$p_1 = s_1 \quad (17)$$

$$p_2 = -s_1 + s_2 \quad (18)$$

$$p_3 = s_1 - s_2 + s_3 \quad (19)$$

$$p_4 = -s_1 + s_2 - s_3 + s_4 \quad (20)$$

$$p_5 = s_1 - s_2 + s_3 - s_4 + s_5 \quad (21)$$

etc.

Several examples of equivalences between topological indices have been reported in the literature [10,11]: it was found that the well-known Wiener index [12] and the Schultz index [13] are related and the resistance distance [14] and the “quasi Wiener index” [15] are also related. Basak et al. [16] after inspecting more than 200 topological indices concluded that the indices could be grouped into various (more or less distinct) classes. Klein and Gutman [17] found that distance based indices are all mutually related and can be expressed in terms of distance distribution moments. Finally we have shown that s_2 is equivalent with one of the Zagreb indices [7].

As a matter of fact it is questionable whether a graph-theoretical index exists, which is *not* related to any other (graph) invariant at all. Skvortsova et al. [18] found several invariants they considered basic. A similar result was obtained by Estrada and Rodriguez [19], who found that the edge-connectivity index does not (practically) depend on other considered types of topological indices.

Path-counts p_{k-1} , p_k and the graph valence shells s_k are linear transforms, therefore, and in contradiction to Randić’s conclusion, they *do not* represent independent features. On the other hand, if such (qualitatively different) indices like path counts and graph valence shells are equivalent, then - using Kirby’s wording [20] - what does a topology index index?

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